

# CSC 495/583 Topics of Software Security

## X86 Assembly & Stack & Stack Frame

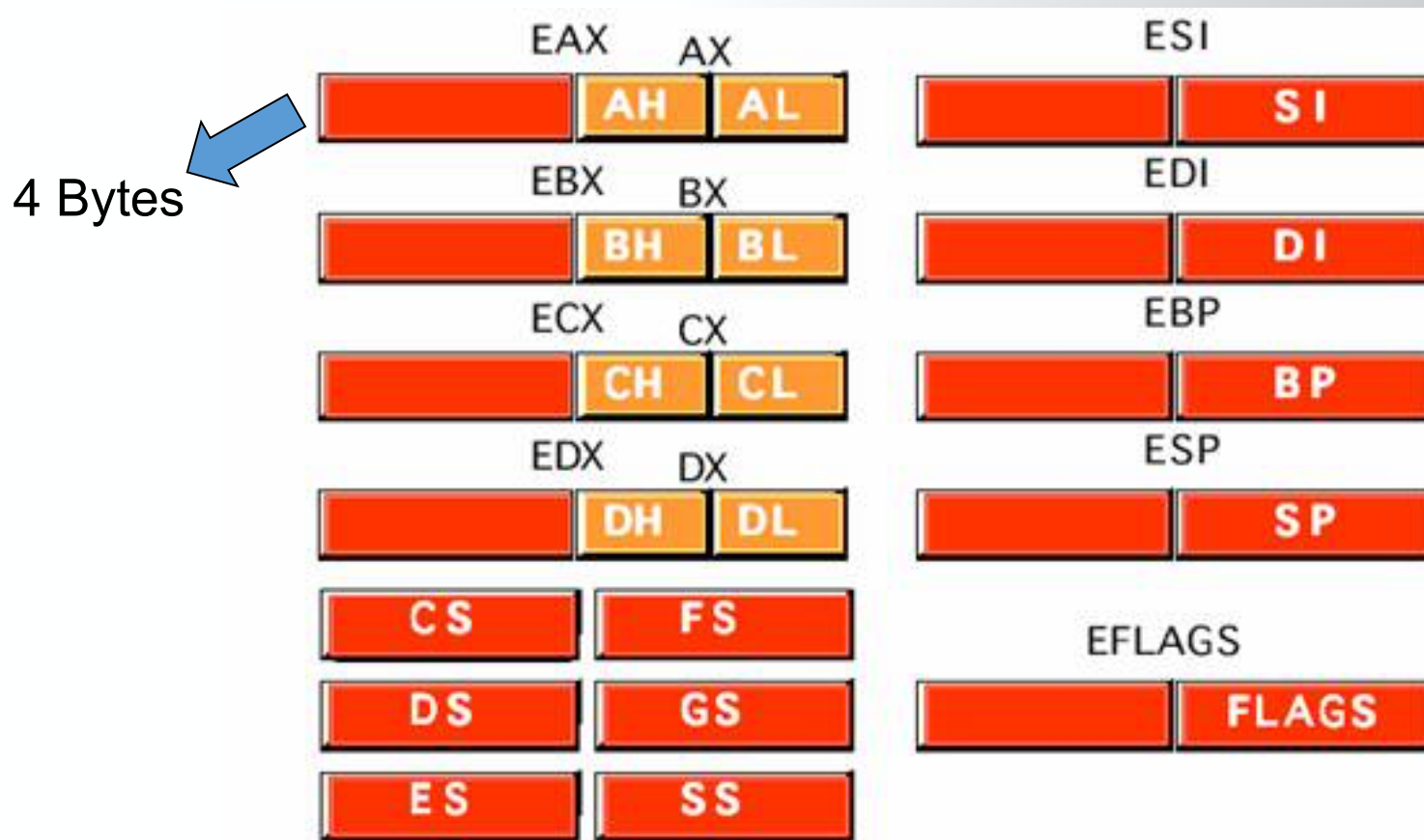
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# Review

# General-purpose Registers

- The **eight** 32-bit general-purpose data registers are used to hold operands for logical and arithmetic operations, operands for address calculations and memory pointers



# X86 ASM

# MOV

- Move **reg/mem** value to **reg/mem**

- mov A, B is "Move B to A" (A=B)
- Same data size

**mov eax, 0x1337**

**mov bx, ax**

**mov [esp+4], bl**

**0x1234**

# More About Memory Access

- `mov ebx, [esp + eax * 4]` **Intel**
- `mov (%esp, %eax, 4), %ebx` **AT&T**
- `mov BYTE [eax], 0x0f`

You must indicate the data size: BYTE/WORD/DWORD

# ADD / SUB

- ADD / SUB
- Normally "reg += reg" or "reg += imm"
- Data size should be equal
  - add eax, ebx
  - sub eax, 123
  - sub eax, BL ; Illegal

# Jump

- Unconditional jump: jmp
- Conditional jump: je/jne and ja/jae/jb/jbe/jg/jge/jl/jle ...
- Sometime with "cmp A, B" -- compare these two values and set eflags
- Conditional jump is decided by some of the eflags bits.

## The JMP Instruction

- JMP (jump) instruction causes an unconditional jump
- Syntax is: **JMP destination/target\_label**
- JMP can be used to get around the range restriction [126/127 byte]
- Flags – no change

TOP:

```
; body of the loop, say 2 instructions
DEC  CX      ; decrement counter
JNZ  TOP     ; keep looping if CX > 0
MOV  AX, BX
```

TOP:

```
; the loop body contains so many instructions
; that label TOP is out of range for JNZ. Solution is-
      DEC  CX
      JNZ  BOTTOM
      JMP  EXIT
```

BOTTOM:

```
JMP  TOP
```

EXIT:

```
MOV  AX, BX
```

Section 6-3: Assembly Language Programming

## Unsigned and Signed Jumps.

Condition	Unsigned	Signed
source < dst	JB	JL
source <= dst	JBE	JLE
source != dst	JNE(JNZ)	JNE(JNZ)
source = dst	JE(JZ)	JE(JZ)
source >= dst	JAE	JGE
source > dst	JA	JG



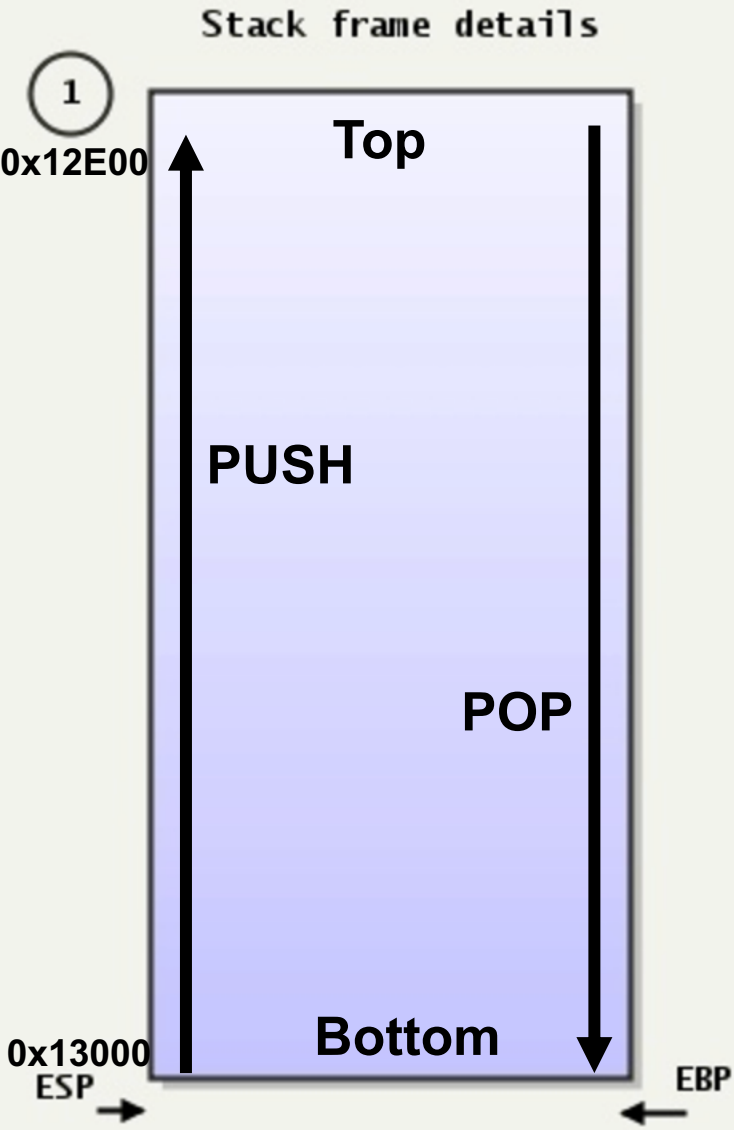
# Jump

- ja/jae/jb/jbe are unsigned comparison
- jg/jge/jl/jle are signed comparison

Unsigned and Signed Jumps.

Condition	Unsigned	Signed
$\text{source} < \text{dest}$	JB	JL
$\text{source} \leq \text{dest}$	JBE	JLE
$\text{source} \neq \text{dest}$	JNE(JNZ)	JNE(JNZ)
$\text{source} = \text{dest}$	JE(JZ)	JE(JZ)
$\text{source} \geq \text{dest}$	JAЕ	JGE
$\text{source} > \text{dest}$	JA	JG

# The Stack



## Stack:

- A special region of your computer's memory that **stores temporary variables** created by each functions
- The stack is a "**LIFO**" (last in, first out) data structure
- Once a stack variable is freed, that region of memory becomes available for other stack variables.

## Properties:

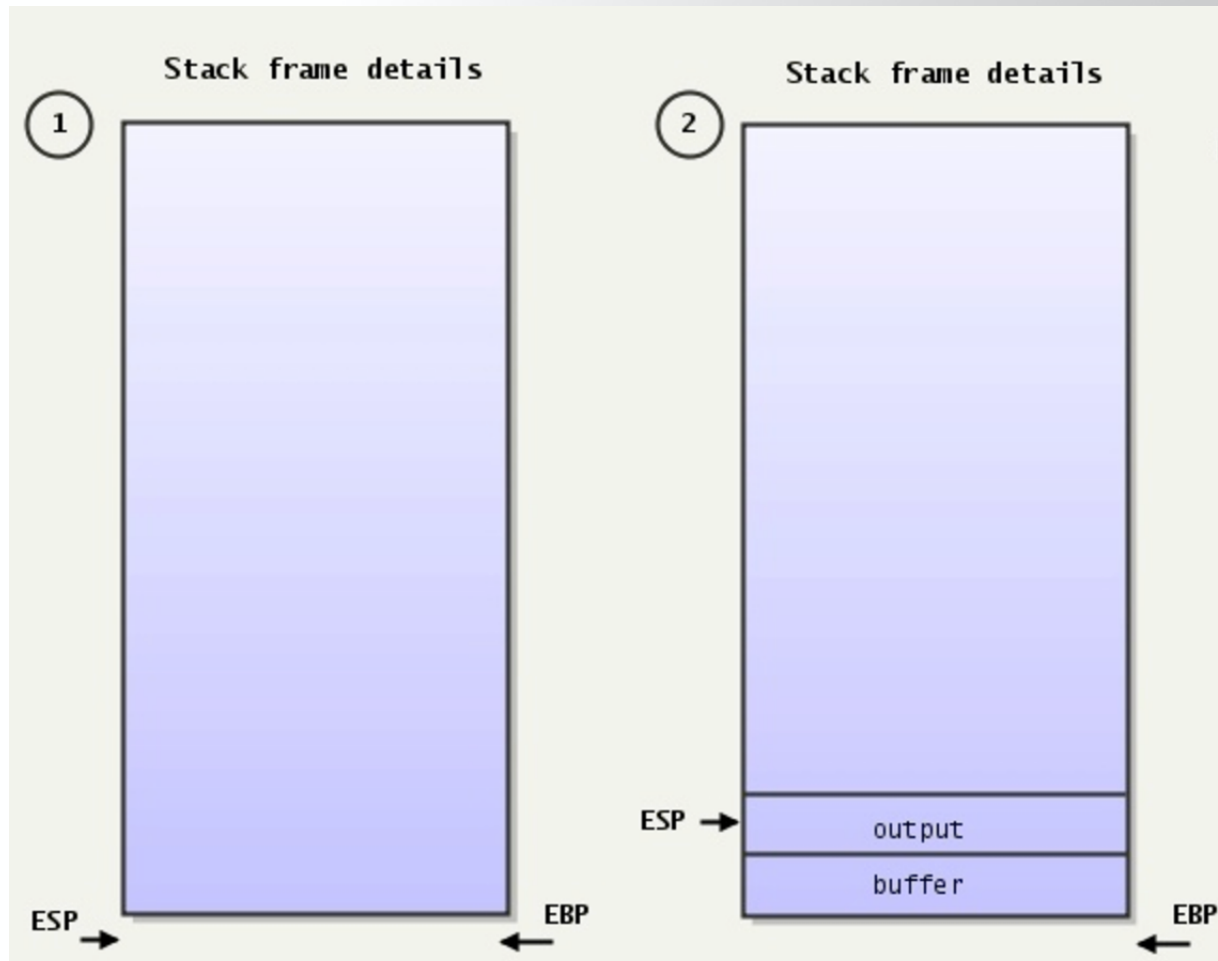
- the stack grows and shrinks as functions **push and pop** local **variables**
- there is no need to manage the memory yourself, variables are allocated and freed **automatically**
- the **stack has size limits**
- stack variables only exist while the function that created them, is running

**EBP—Pointer to data on the stack**  
**ESP—Stack pointer**

# The Stack

## Stack:

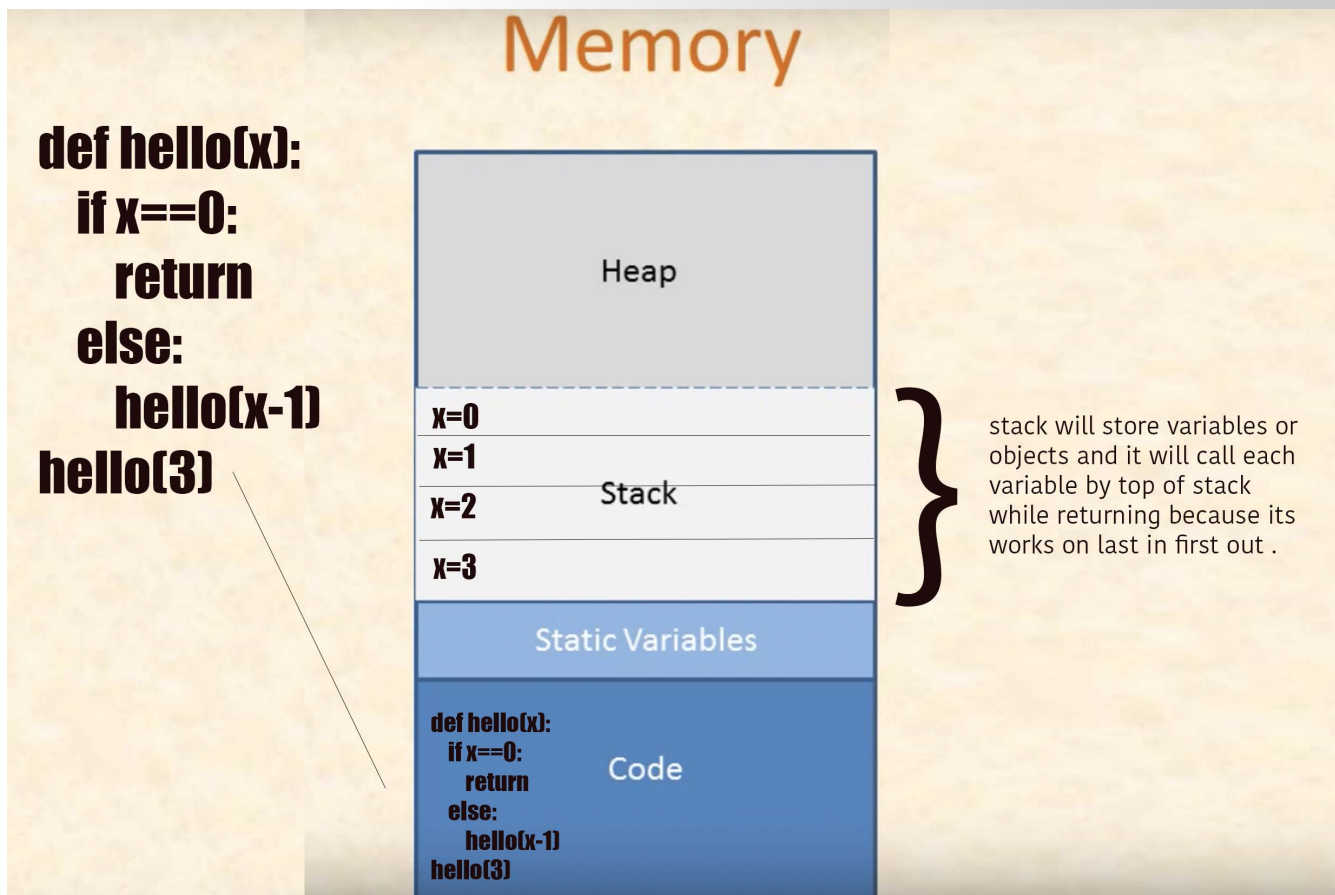
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# Stack Frame

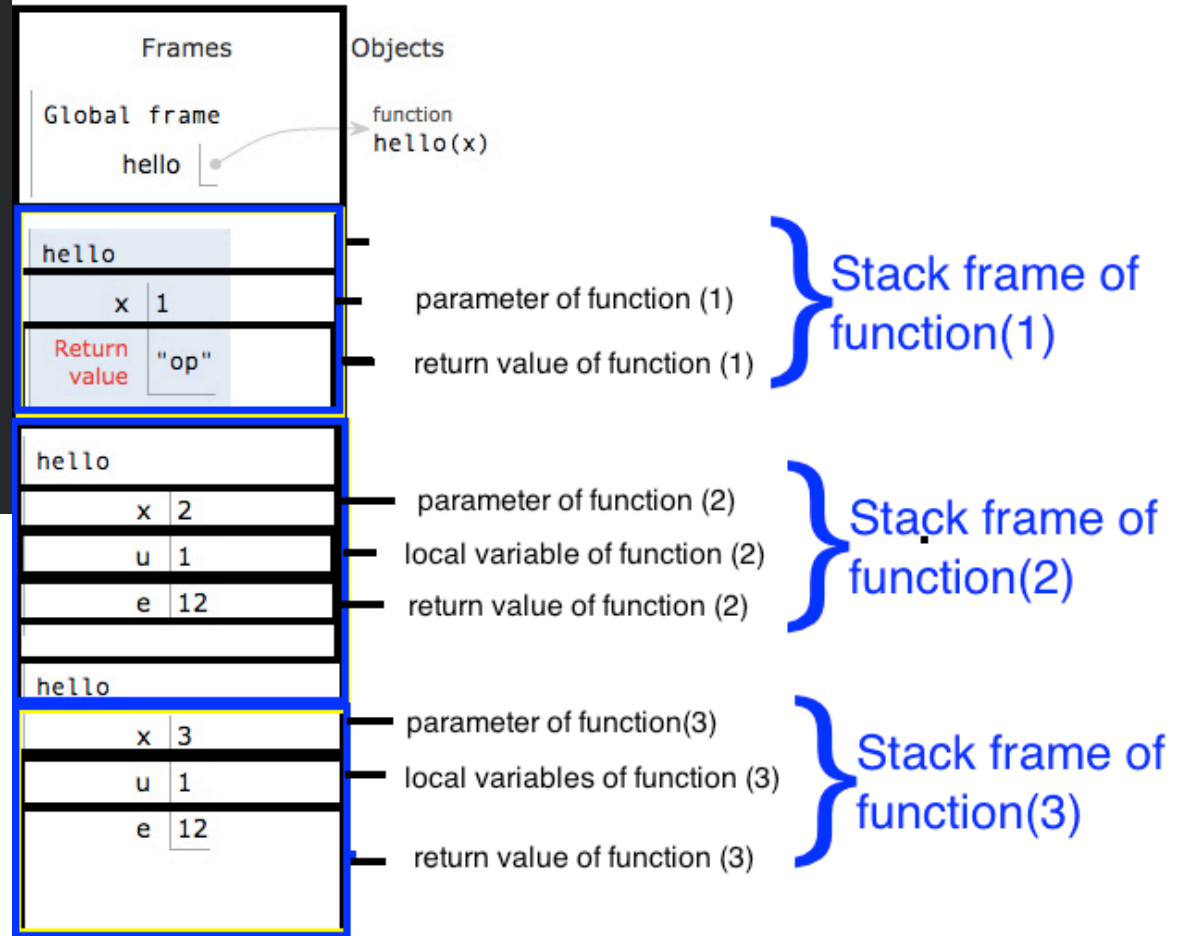
# Stack Frame

- A stack frame is **a frame of data that gets pushed onto the stack.**
- In the case of a **call stack**, a stack frame would represent **a function call and its argument data.**



# Stack Frame

```
1 def hello(x):
2     if x == 1:
3         return "op"
4     else:
5         u = 1
6         e = 12
7         s = hello(x - 1)
8         e += 1
9         print(s)
10        print(x)
11        u += 1
12    return e
13
14
15 hello(3)
```



# Functions and Frames

Each function call results in a new frame being created on the stack.

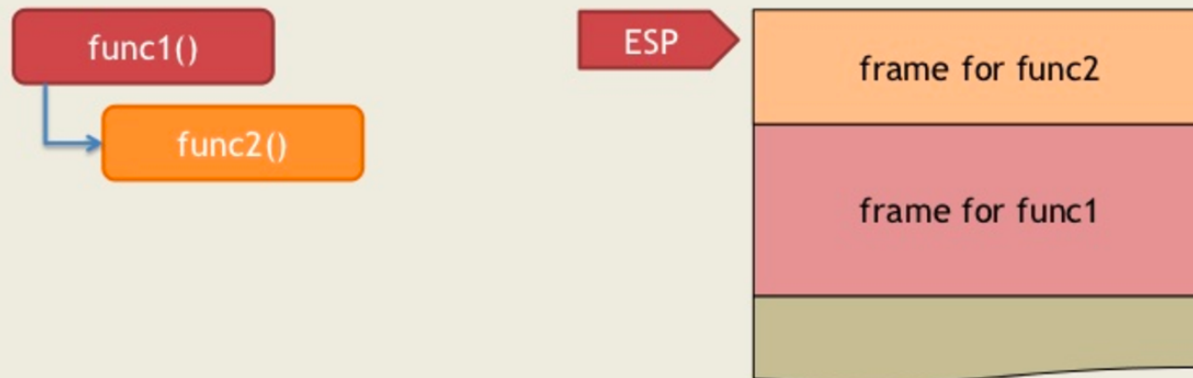
func1()

ESP

frame for func1

# Functions and Frames

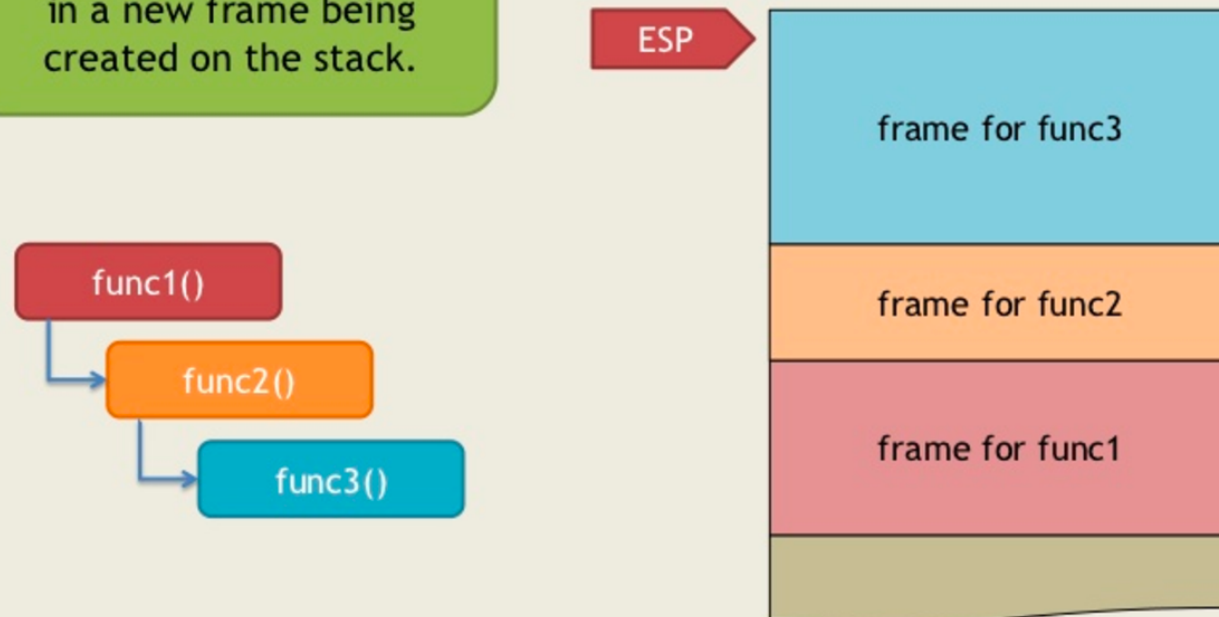
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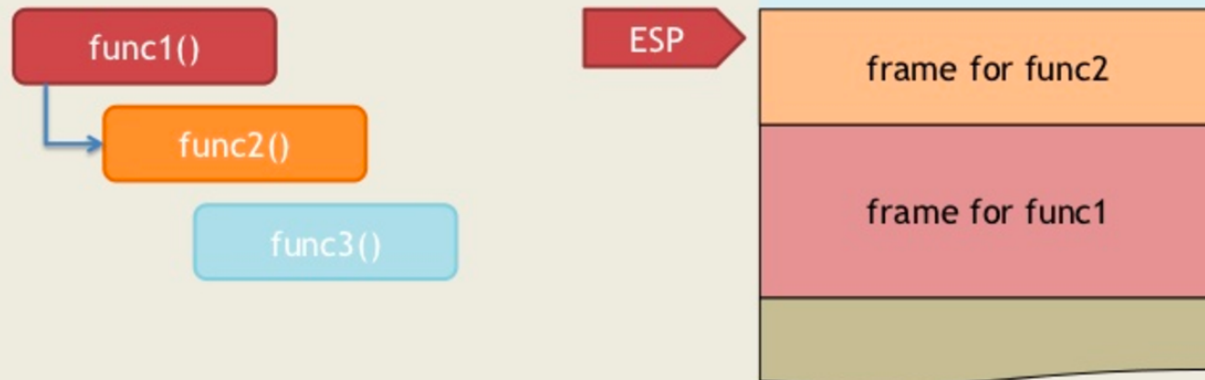
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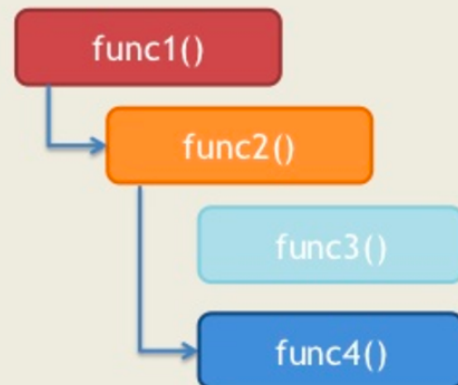
# Functions and Frames

When a function returns, the frame is "unwound" or "collapsed".

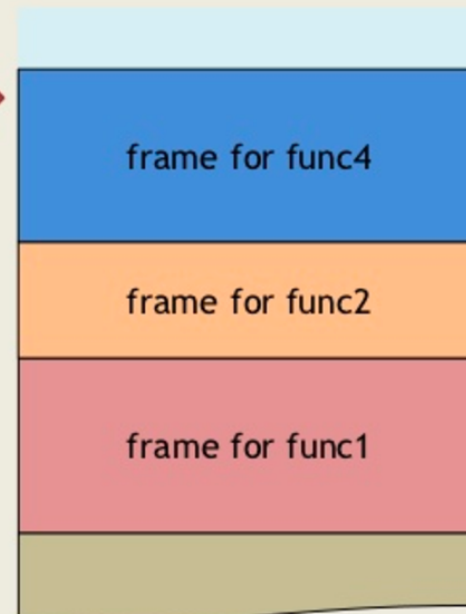


# Functions and Frames

And as new functions get invoked, new frames get created.



ESP



# Stack Frame

File Edit View Terminal Tabs Help

```
PUSH EBP      ; start of the func (save current EBP to stack)
MOV EBP, ESP  ; save current ESP to EBP

.....      ; function body
              ; no matter how ESP changes, the EBP remains unchanged

MOV ESP, EBP  ; move the saved function start addr back to ESP
POP EBP       ; before return the func, pop the stored EBP
RETN          ; end of the func
```

-- INSERT --

12,1

All

# StackFrame.c

```
1 StackFrame.c +
  1 #include "stdio.h"
  2
  3 long add(long a, long b)
  4 {
  5     long x = a, y = b;
  6     return (x + y);
  7 }
  8
  9 int main(int argc, char* argv[])
10 {
11     long a = 1, b = 2;
12     printf("%d\n", add(a,b));
13     return 0;
14 }
15
```

# Q & A

