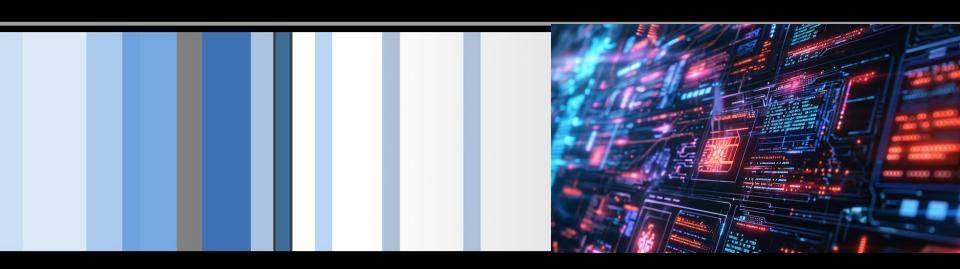
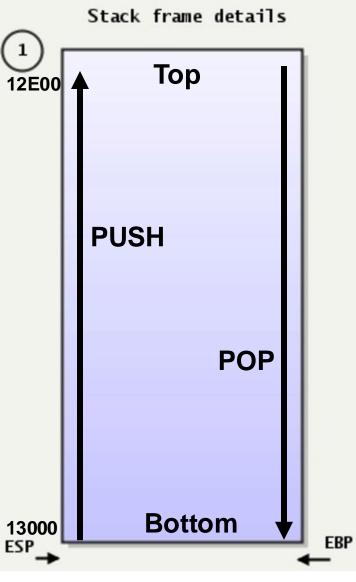


CSC 600 Advanced Seminar Stack Frame & Calling Convention

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The Stack



Stack:

- A special region of your computer's memory that **stores temporary variables** created by each functions
- The stack is a "LIFO" (last in, first out) data structure
- Once a stack variable is freed, that region of memory becomes available for other stack variables.

Properties:

- the stack grows and shrinks as functions push and pop local variables
- there is no need to manage the memory yourself, variables are allocated and freed automatically
- the stack has size limits
- stack variables only exist while the function that created them, is running

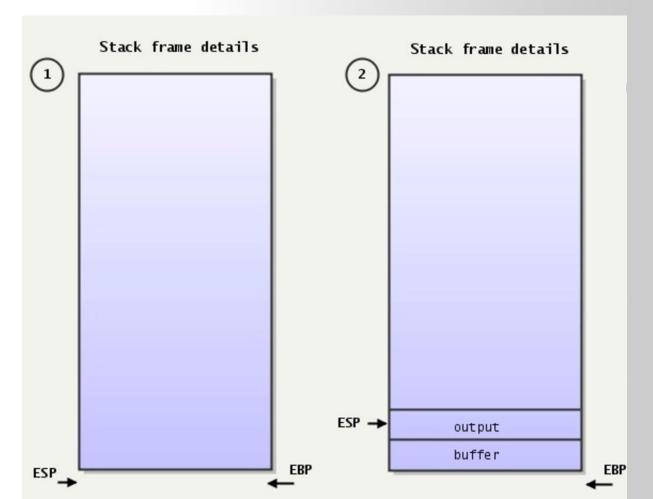
EBP—Pointer to data on the stack ESP—Stack pointer



The Stack

Stack:

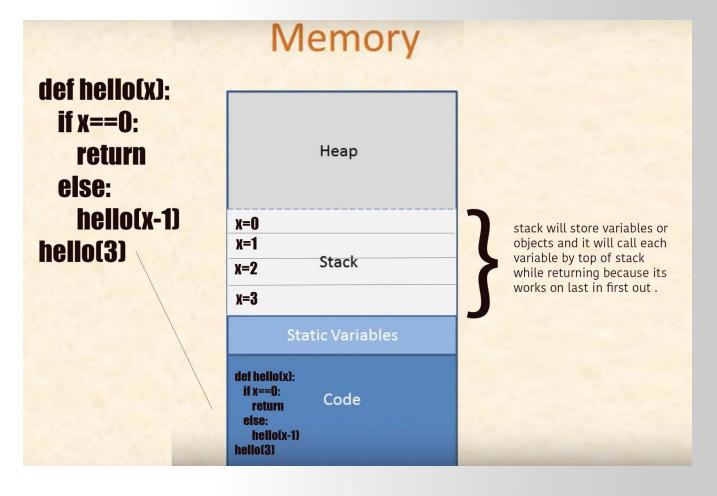
- A special region of your computer's memory that stores temporary variables created by each functions
- The stack is a "LIFO" (last in, first out) data structure
- Once a stack variable is freed, that region of memory becomes available for other stack variables.



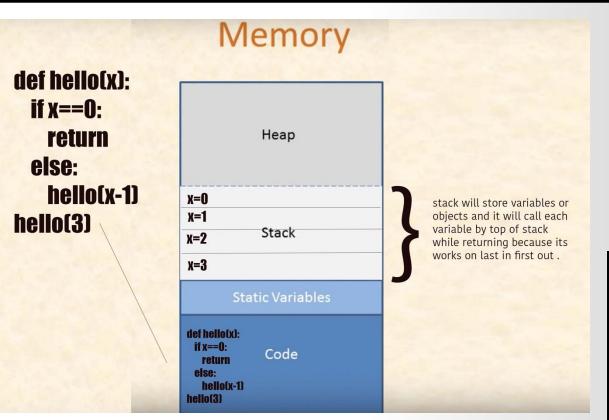




- A stack frame is a frame of data that gets pushed onto the stack.
- In the case of a call stack, a stack frame would represent a function call and its argument data.







```
1 def hello(x):
2     if x == 0:
3         return
4     else:
5         hello(x-1)
6
7 hello(9999999)
```

```
hello(x-1)
  File "stack.py", line 5, in hello
    hello(x-1)
  File "stack.py", line 5, in hello
   hello(x-1)
  File "stack.py", line 5, in hello
    hello(x-1)
  File "stack.py", line 5, in hello
    hello(x-1)
  File "stack.py", line 5, in hello
    hello(x-1)
RuntimeError: maximum recursion depth exceeded
 quake0day@quakes-iMac
```

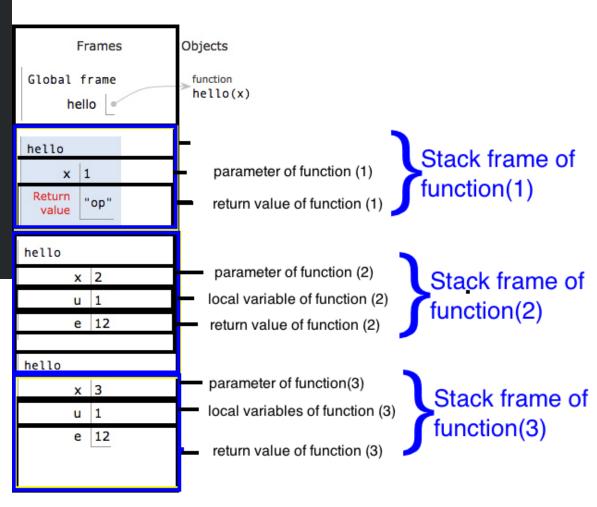
Stack

- Pass arguments
- Save the return address
- Save local variable



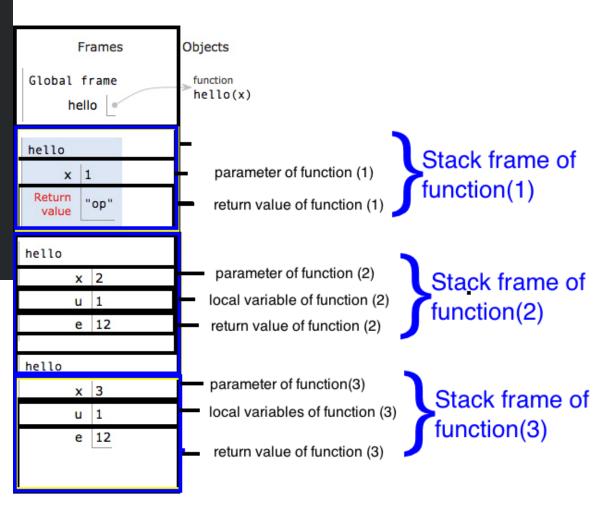
```
def hello(x):
    if x == 1:
        return "op"
    else:
        u = 1
        e = 12
        s = hello(x - 1)
        e += 1
        print(s)
        print(x)
        u += 1
    return e

hello(3)
```



```
def hello(x):
    if x == 1:
        return "op"
    else:
        u = 1
        e = 12
        s = hello(x - 1)
        e += 1
        print(s)
        print(x)
        u += 1
    return e

hello(3)
```



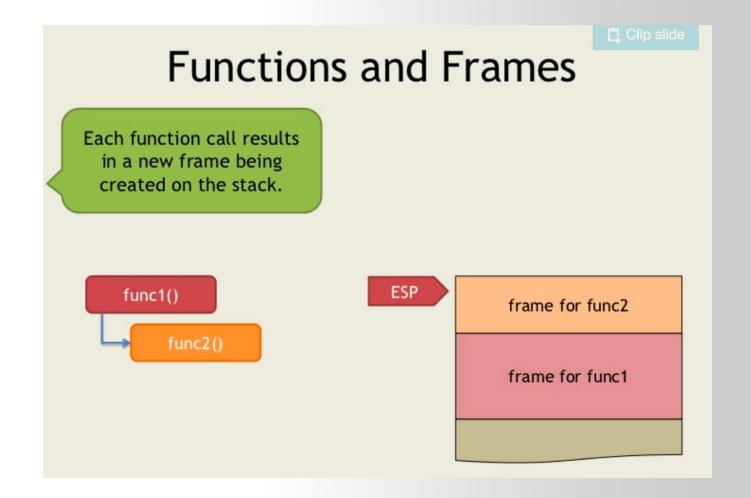
Local Variable

- Limited Register(s) → Store *Local Variable* in stack
 - Use esp and ebp to define a stack frame for current function
 - Use relative position of esp or ebp for retrieving and storing data
 - e.g. mov eax, [esp+124]
- Very easy to do recursive call

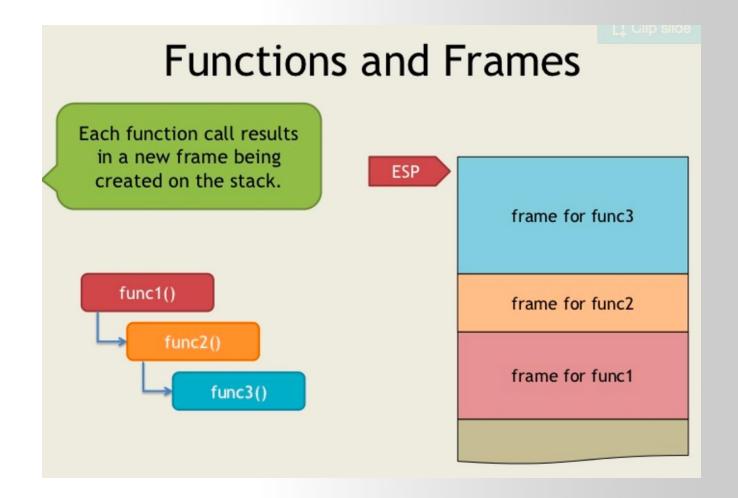


Functions and Frames Each function call results in a new frame being created on the stack. func1() **ESP** frame for func1

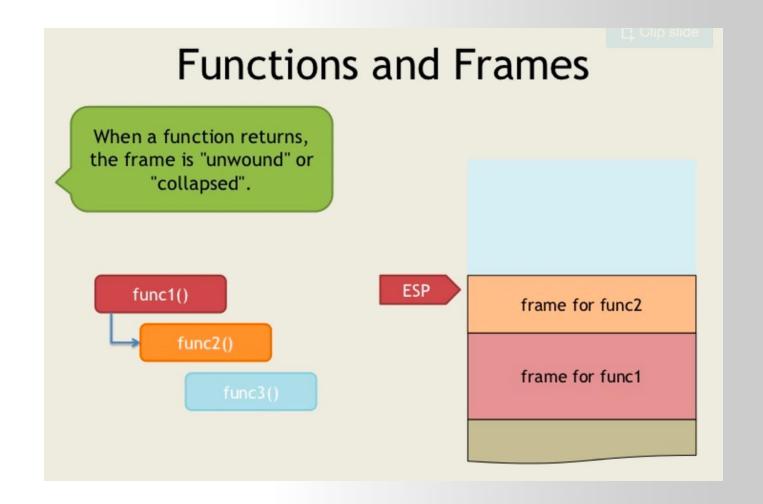














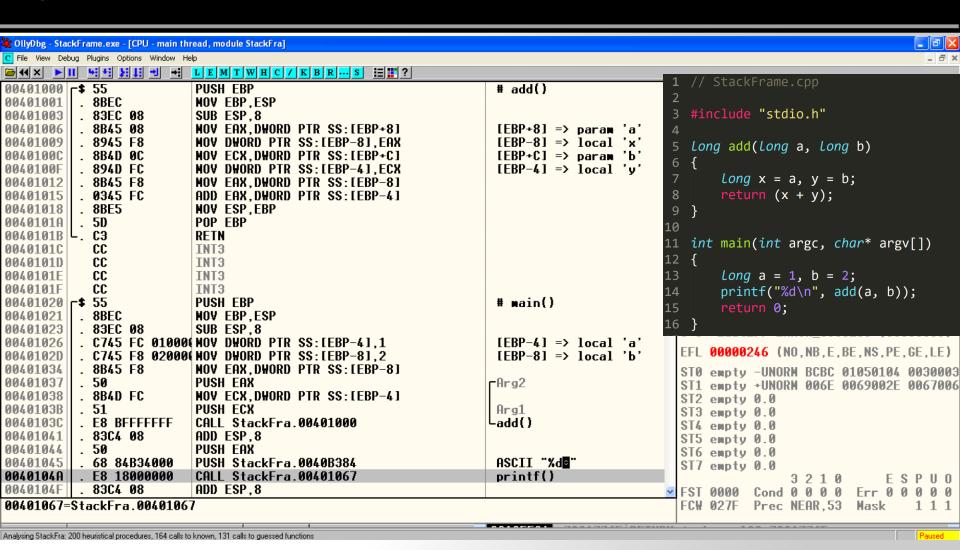
Functions and Frames And as new functions get invoked, new frames get created. **ESP** frame for func4 func1() frame for func2 frame for func1 func4()



```
File Edit View Terminal Tabs Help
PUSH EBP ; start of the func (save current EBP to stack)
MOV EBP, ESP ; save current ESP to EBP
               ; function body
               ; no matter how ESP changes, the EBP remains unchanged
MOV ESP, EBP ; move the saved function start addr back to ESP
        ; before return the func, pop the stored EBP
POP EBP
RETN
               ; end of the func
  INSERT --
                                                             12,1
                                                                           All
```

```
// StackFrame.cpp
 2
   #include "stdio.h"
 4
   Long add(Long a, Long b)
 6
       long x = a, y = b;
       return (x + y);
 8
 9
10
   int main(int argc, char* argv[])
12
13
       long a = 1, b = 2;
       printf("%d\n", add(a, b));
14
15
       return 0;
16 }
```



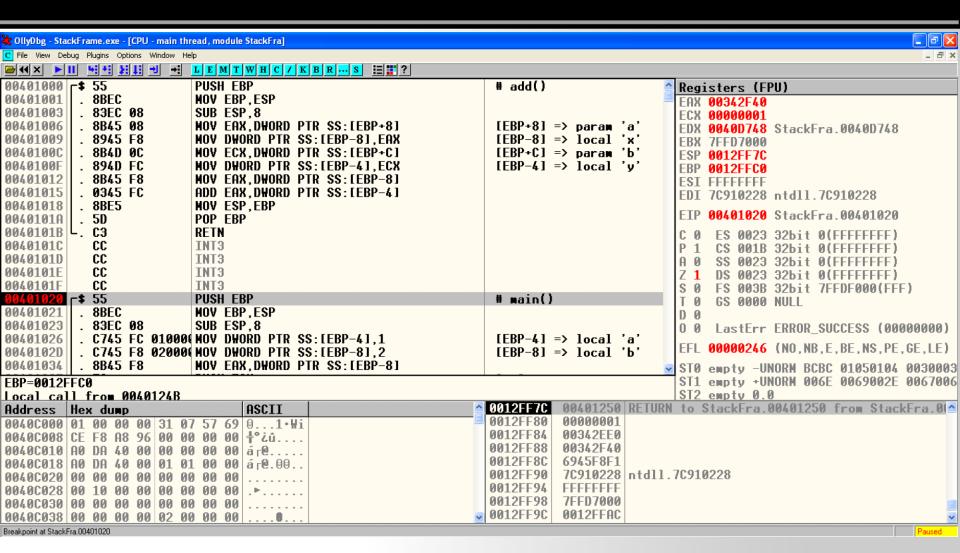


EBP - n: Local vars EBP + n: Parameters

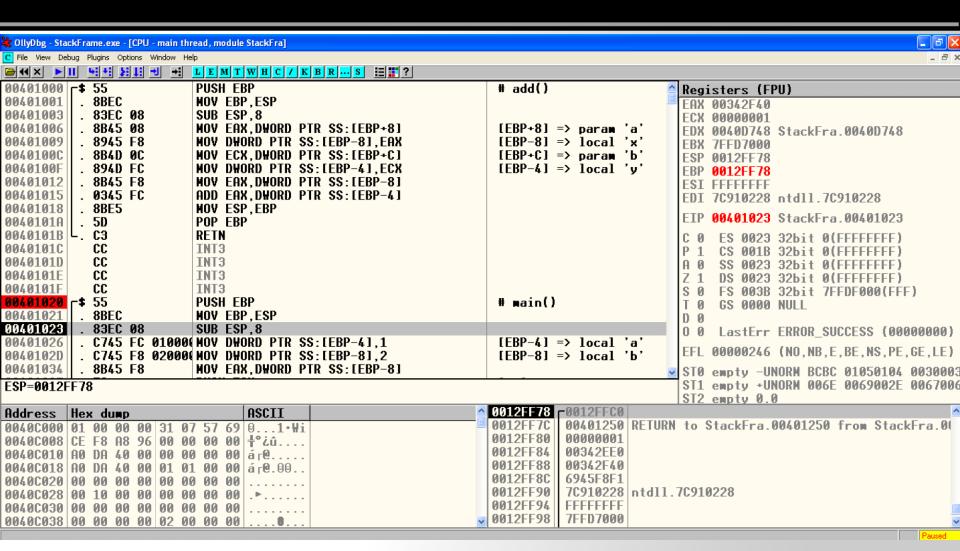


```
// StackFrame.cpp
 3
   #include "stdio.h"
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   long add(Long a, Long b)
 6
        long x = a, y = b;
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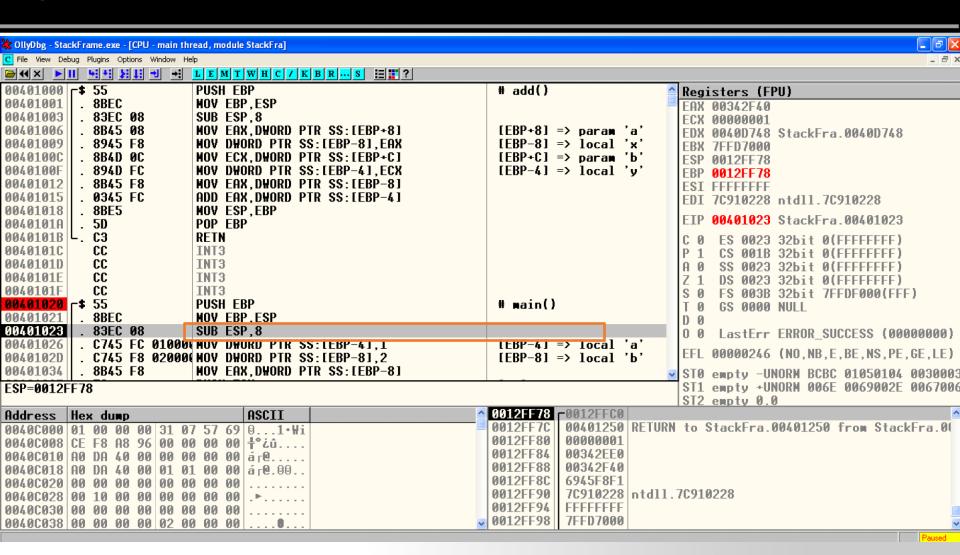






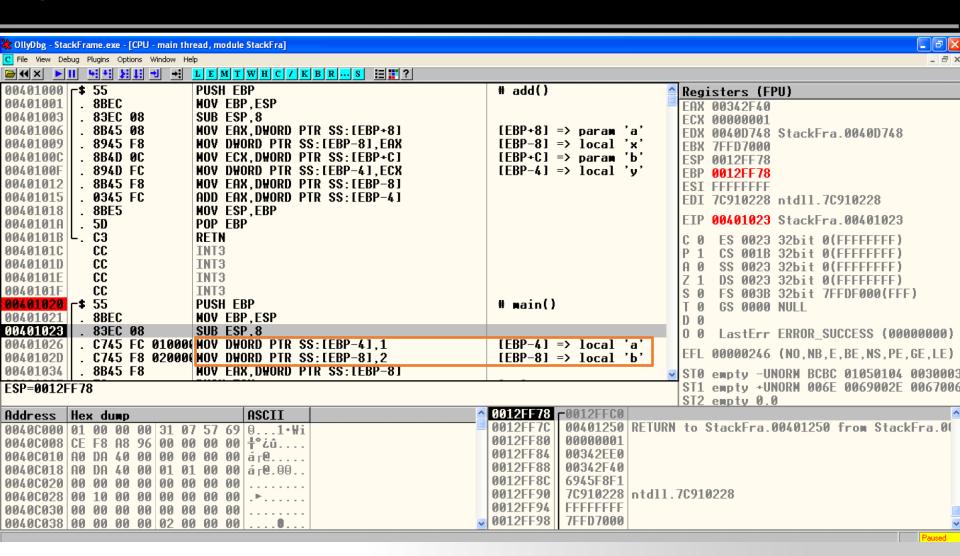
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Create space for 'a' and 'b' → long → 4 byte



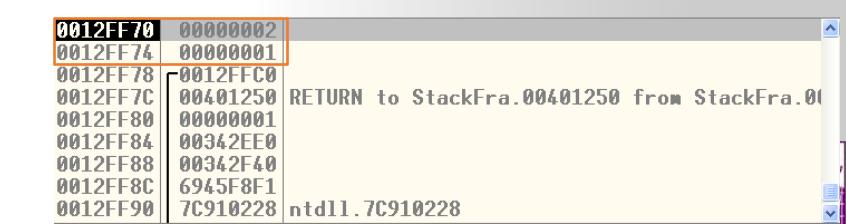


Create space for 'a' and 'b' → long → 4 byte



Assembly	C	Type Conversion
DWORD PTR SS:[EBP-4]	*(DWORD*)(EBP-4)	DWORD (4 byte)
WORD PTR SS:[EBP-4]	*(WORD*)(EBP-4)	WORD (2 byte)
BYTE PTR SS:[EBP-4]	*(BYTE*)(EBP-4)	1 byte

4 Byte memory space at address [EBP-4]



```
1 // StackFrame.cpp
   #include "stdio.h"
4
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       long x = a, y = b;
       return (x + y);
8
10
   int main(int argc, char* argv[])
11
12
13
       long a = 1, b = 2;
       printf("%d\n", add(a, b));
14
       return 0;
15
16 }
```

00401034	. 8B45 F8	MOV EAX, DWORD PTR SS:[EBP-8]	
00401037	. 50	PUSH EAX	rArg2
00401038	. 8B4D FC	MOV ECX, DWORD PTR SS:[EBP-4]	
0040103B	. 51	PUSH ECX	Arg1 -add()
0040103C	. 50 . 8B4D FC . 51 . E8 BFFFFFF	CALL StackFra.00401000	Ladd()



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```

00401000	┌\$ 55	PUSH EBP	# add()
00401001	. 8BEC	MOV EBP, ESP	



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```

00401000 -\$	55	PUSH EBP	# add()
00401001		MOV EBP, ESP	
00401003 . 00401006 .	83EC 08	SUB ESP,8	
00401006	8B45 08	MOV EAX,DWORD PTR SS:[EBP+8]	[EBP+8] => param 'a'
00401009 .	8945 F8	MOV DWORD PTR SS:[EBP-8],EAX	[EBP-8] => local 'x'
0040100C .	8B4D 0C	MOV ECX,DWORD PTR SS:[EBP+C]	[EBP+C] => param 'b'
0040100C . 0040100F .	894D FC	MOV DWORD PTR SS:[EBP-4],ECX	[EBP-4] => local 'y'

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       return 0;
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```

```
00401018 | . 8BE5 | MOV ESP,EBP | POP EBP | RETN
```



	. 83C4 08 . 50	ADD ESP,8 Clean Stack	
00401045	. 68 84B34000	PUSH StackFra.0040B384	ASCII "%de"
0040104A		CALL_StackFra.00401067	printf()



00401041 00401044	. 83C4 08 . 50	ADD ESP,8 Clean Stack	
00401045	. 68 84B34000	PUSH StackFra.0040B384	ASCII "%de"
0040104A		CALL_StackFra.00401067	printf()



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       return 0;
16
```

```
00401052
00401054
00401056
```

. 33C0 . 8BE5 . 5D XOR EAX,EAX MOV ESP,EBP POP EBP Set EAX -> 0
Faster than
MOV EAX,0



Calling Convention



Two Questions

Q: When a function finished, how to handle the parameter left in the stack.

0012FF70	00000002		^
0012FF74	00000001		
0012FF78	┌0012FFC0		
0012FF7C	00401250	RETURN to StackFra.00401250 from StackFra.0)(
0012FF80	00000001		
0012FF84	00342EE0		
0012FF88	00342F40		
0012FF8C	6945F8F1		
0012FF90	70910228	ntdl1.7C910228	~

A: We don't care...

Q: When a function finished, how change the ESP value?

A: ESP should be restored to the previous value



Standard C Calling Conventions

- Calling conventions are a standardized method for functions to be implemented and called by the machine.
- A calling convention specifies the method that a compiler sets up to access a subroutine.
- There are three major calling conventions that are used with the C language on 32-bit x86 processors:
 - CDECL
 - STDCALL,
 - FASTCALL.



CDECL

- The C language, by default, uses the CDECL calling convention
- In the CDECL calling convention the following holds:
 - Arguments are passed on the stack in Right-to-Left order, and return values are passed in eax.
 - The calling function cleans the stack. This allows CDECL functions to have variable-length argument lists.



STDCALL

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- In the CDECL calling convention the following holds:
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 The calling function cleans the stack. This allows CDECL functions to have variable-length argument lists.

```
_cdecl int MyFunction1(int a, int b)
    return a + b;
and the following function call:
  x = MyFunction1(2, 3);
These would produce the following assembly listings, respectively:
 _MyFunction1:
 push ebp
 mov ebp, esp
 mov eax, [ebp + 8]
 mov edx, [ebp + 12]
 add eax, edx
 pop ebp
 ret
and
 push 3
 push 2
 call MyFunction1
  add esp, 8
```

STDCALL

- STDCALL, also known as "WINAPI" (and a few other names, depending on where you are reading it) is used almost exclusively by Microsoft as the standard calling convention for the Win32 API.
 - STDCALL passes arguments right-to-left, and returns the value in eax.
 - The called function cleans the stack, unlike CDECL. This means that STDCALL doesn't allow variable-length argument lists.



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 - STDCALL passes arguments right-to-left, and returns the value in eax.
 - The called function cleans the stack, unlike CDECL. This means that STDCALL doesn't allow variable-length argument lists.

RET 8 → RET + POP 8 Byte

```
Consider the following C function:
  _stdcall int MyFunction2(int a, int b)
     return a + b:
and the calling instruction:
   x = MyFunction2(2, 3);
These will produce the following respective assembly code fragm
 : MyFunction2@8
 push ebp
 mov ebp, esp
 mov eax, [ebp + 8]
 mov edx, [ebp + 12]
 add eax, edx
 pop ebp
and
 push 3
 push 2
 call MyFunction208
```

FASTCALL

- The FASTCALL calling convention is not completely standard across all compilers, so it should be used with caution.
- The calling function most frequently is responsible for cleaning the stack, if needed.





